FUSION-10°

Running NoSQL Natively on Flash Fusion-io SDK

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X

The Future of High Performance Storage?

Everyone seems to agree.







ioMemory

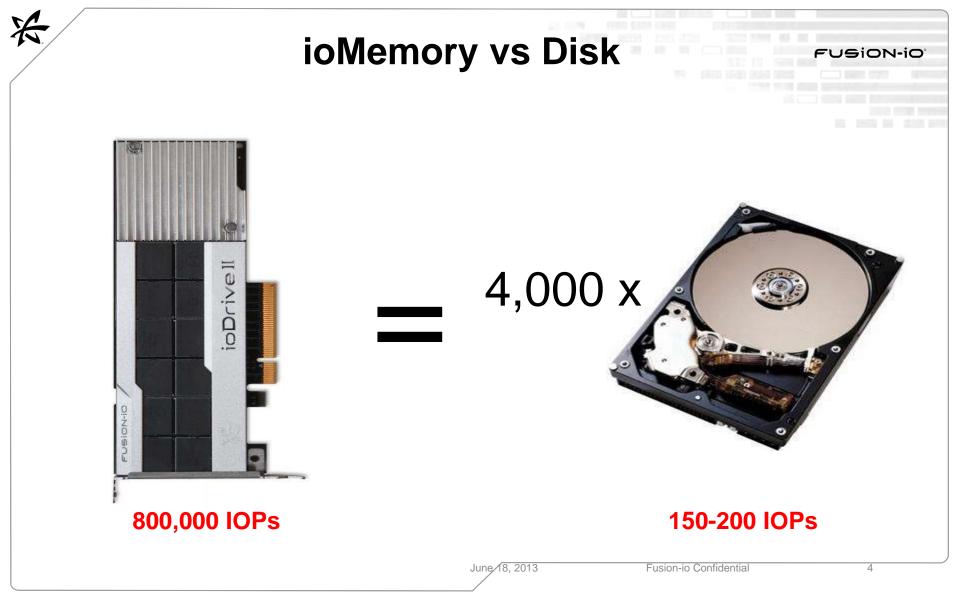
- A New Memory tier called ioMemory
 - Leverages the best advantages of DRAM and rotating drives
 - High Speed like DRAM
 - Persistence and Large capacity of Spinning Hard Drives
- PCIe based NAND Flash storage
 - Micro-second level Disk Access Latency 15µs
 - Very high data throughput 1,5GB/s
 - Very high IOPS 400.000 random write/s
 - Scalable stay ahead of data / performance demanc
 - Advanced wear-leveling algorithm
 - N+1 Chip level redundancy (think RAID protection on card)
 - 100% data integrity protection in case of power loss
 - Endurance is PBW TB's written daily for more than 8 years!

Manufactured by Fusion-io - OEM'ed by









Where to use ioMemory

TIBCO

∮ software[™]

SolidWorks

Ps

NX

NASTRAN

-lustre

IBM GPFS

FUSION-IO

NET.

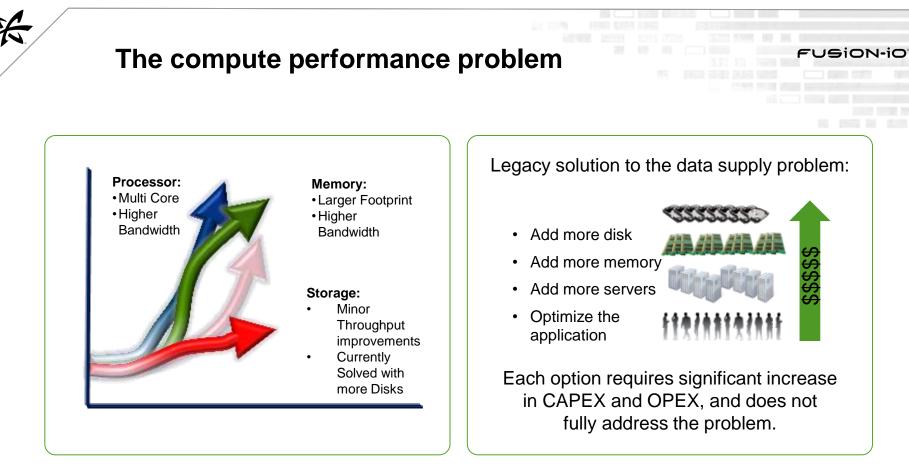


Powered by

Banto

VARNISH

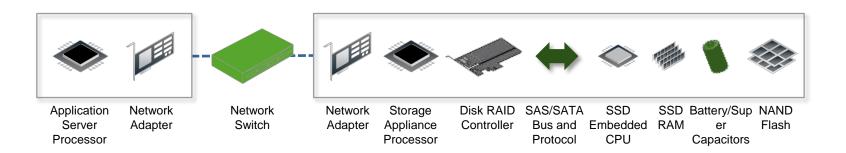
splunk>



"Compute power continues to outpace performance delivered by Storage."

Problem is not getting better, its getting worse."

Networked storage data supply chain from application to flash



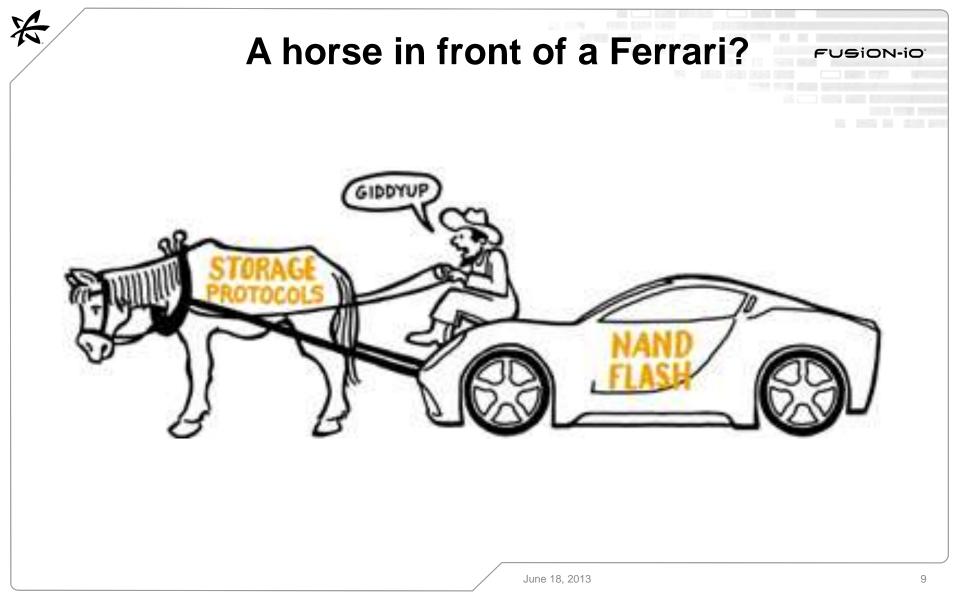
9 Intermediary components required

- All adding access delay, cost, complexity, and lowering reliability (especially the super capacitors)
- Requests must do a round trip touching everything TWICE...

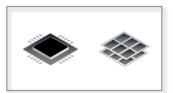
SSD data supply chain from application to flash FUSION-10 Disk RAID SAS/SATA Application SSD SSD Battery/Sup NAND Controller Embedded RAM Server Flash Bus and er CPU Protocol Capacitors Processor

• 5 Intermediary components required

 All adding access delay, cost, complexity, and lowering reliability (especially the super capacitors)



Fusion-io Approach From Application.... to Flash

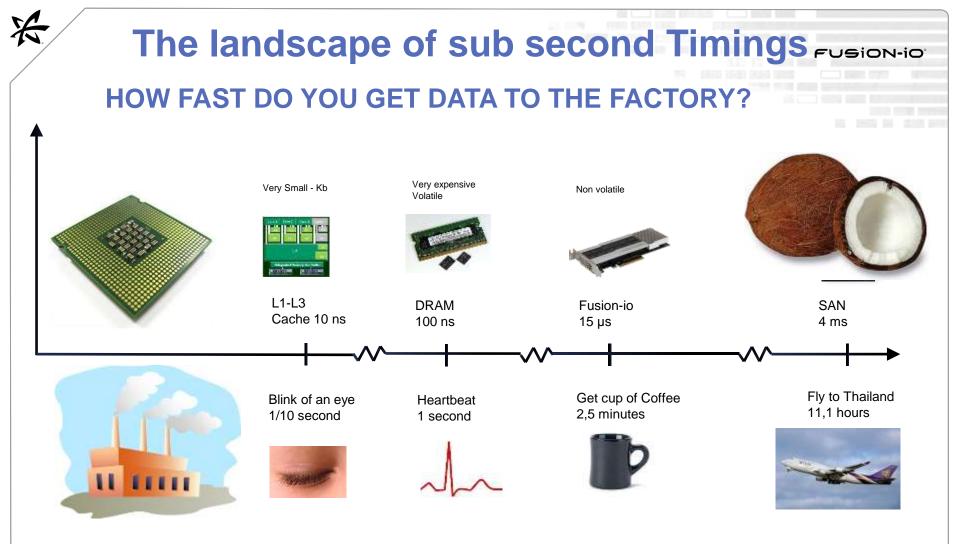


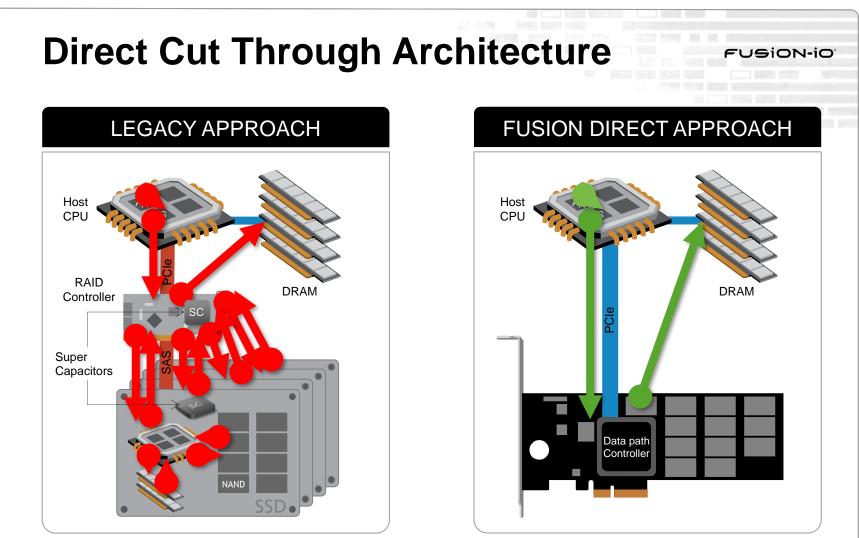
Application N Server F Processor

NAND Flash

• 0 Intermediary components required

No need for super capacitors because data is not "buffered" in DRAM FUSION-iO





The goal of every I/O operation is to move data to and from DRAM and the device

Fusion-io is not a SSD device

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FUSION-iO



Usage Models – Baby Steps

Fusion-io Configential

- Moving specific components of the database to the ioDrives:
 - Tempdb database
 - Indexes
 - Frequently accessed tables
 - Transaction logs-
 - Partition tables







All In

 If database size permits, placing entire database system on Fusion-ioDrives provides maximum performance benefit

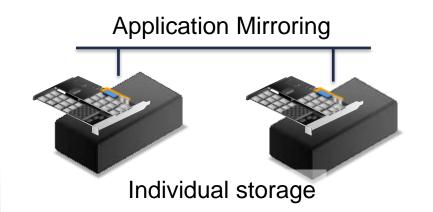






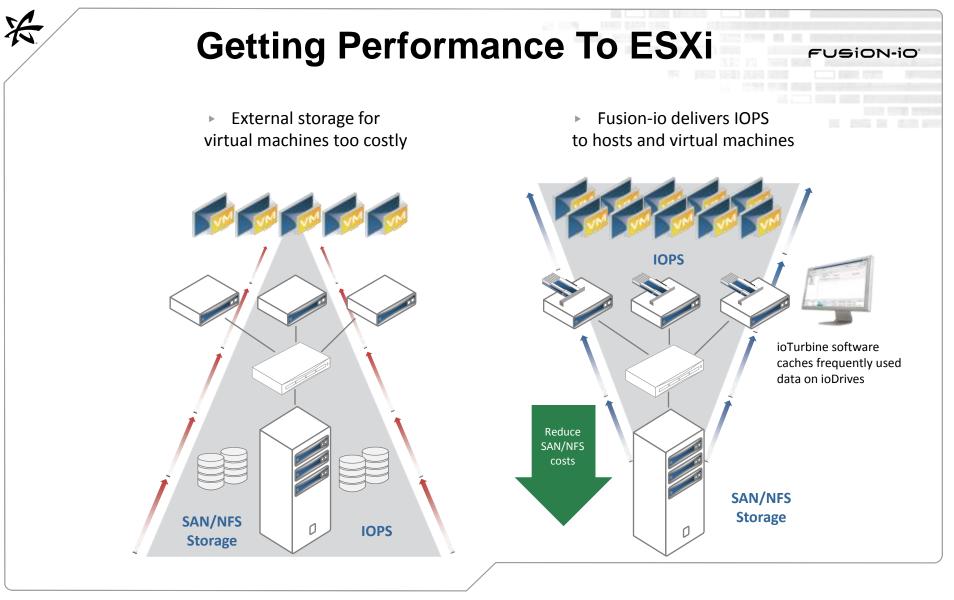
n Node Cluster

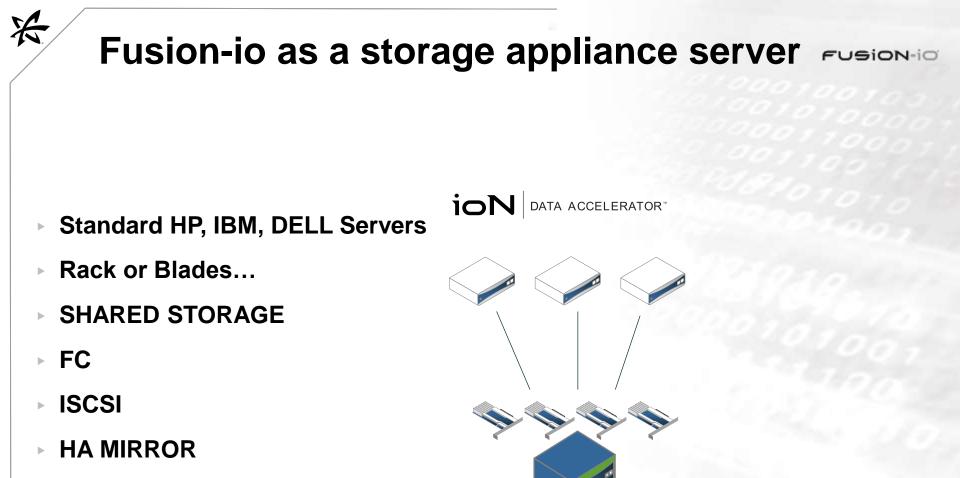
FUSION-iO

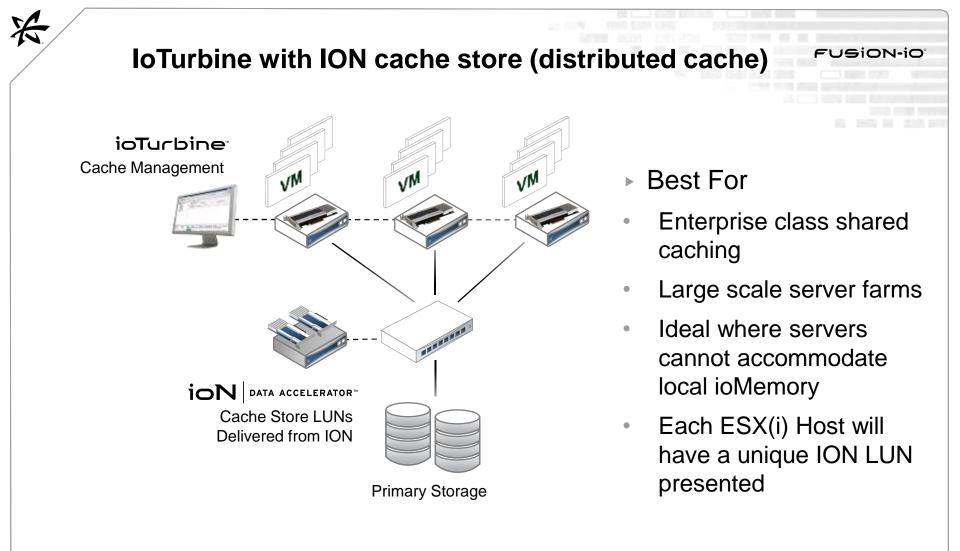


- Perfect NoSQL Model
- MSFT SQL Server Always On
- Oracle Dataguard
- SIOS DataKeeper
- Advantages
 - Fast replication
 - Just another block storage device

Clustering / HA with No Shared Storage!







OS Support

support.fusionio.com

DASHBOARD DOV	WNLOADS	KNOWLEDGE BASE	COMMUNITY	1
1. Identify Product	2.5	lect Files for Download	i -	
Pick Product Options		 IMPORTANT INFO (ou have selected the third ardware into an extensible ncludes: Improved read and writi Greater scalability with Adaptive Flashback⁷¹⁴ t ioMemory® devices to r events 	generation VSL Th platform to fully of e performance multi-CPU archit echnology with so	ecturo elf-hei
Operating System: Linux_rhel-6	Ava	ilable Downloads:		
Version		Documentation		
3.1.1		Driver Source		
		File Name		Siz
		iomemory-vsl-3.1.1.172-	1.0.el6.src.rpm	2.5
		Driver Binaries		
		D Utilities		

C Firmware

D SDK

Windows

Windows Server 2003 SP2 Windows 7 64 bit Windows 8 64 bit (in Oct) Windows Server 2008 R1 SP2 Windows Server 2008 R2 Windows Server 2012 (in Oct)

Linux

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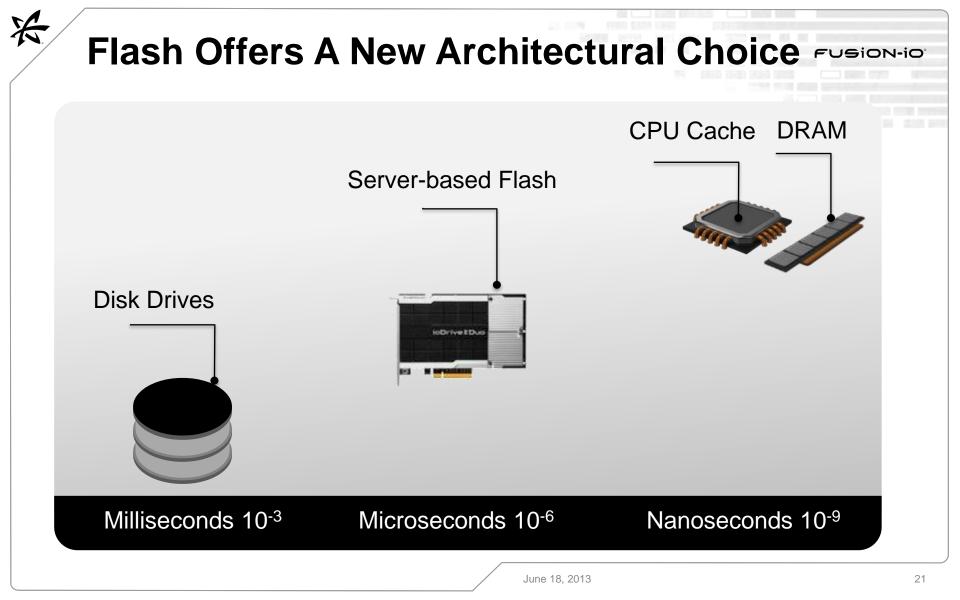
RHEL 5.6, 5.7, 5.8, 6.0, 6.1, 6.2 SLES 10.4, 11, 11.1 OEL 5.7, 6.0, 6.1, 6.2 CentOS 5.6, 5.7, 6.0, 6.1, 6.2 Debian Squeeze Fedora 15, 16 openSUSE 12.1 Ubuntu 10.04, 11.10

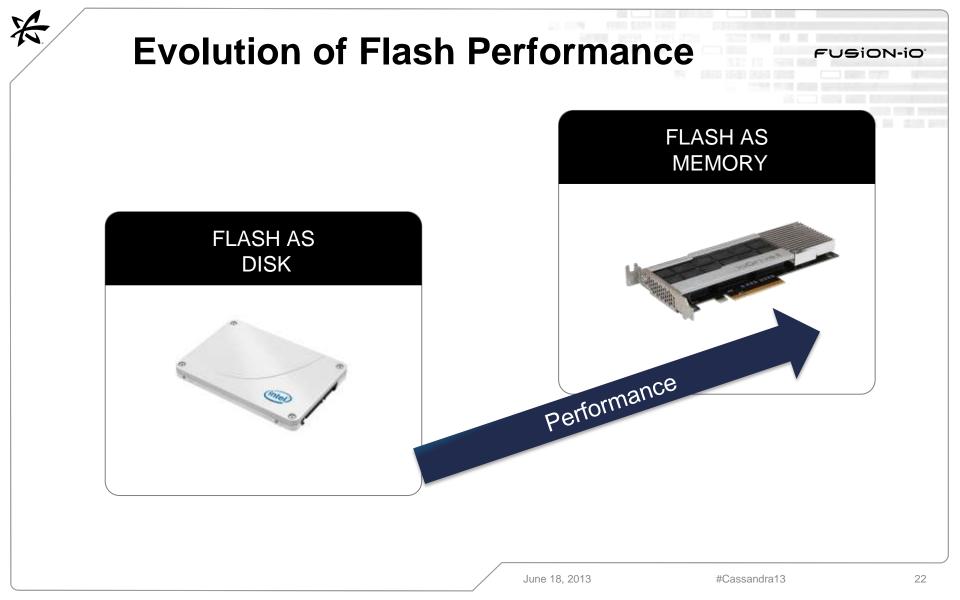
Hypervisors

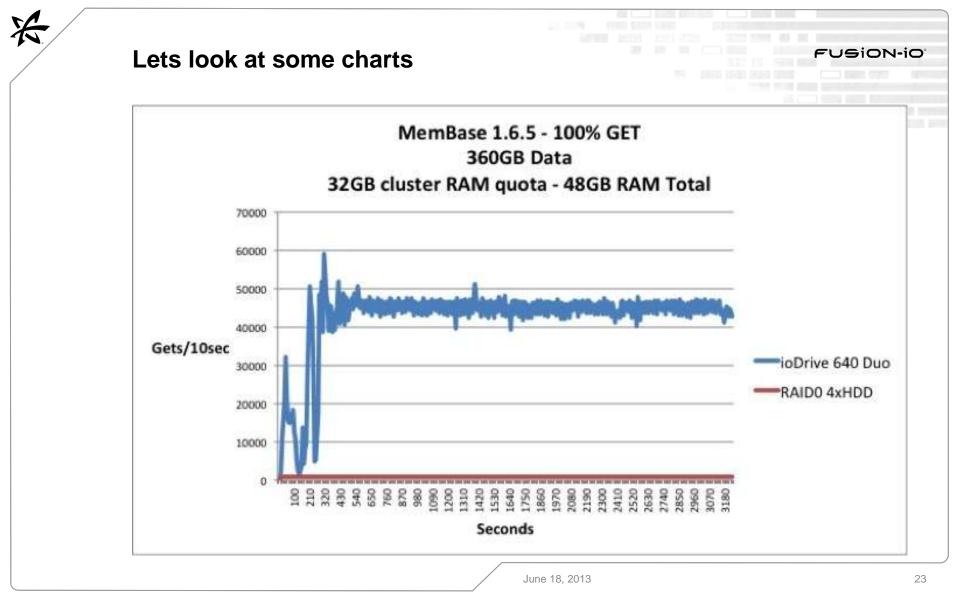
VMware ESX 4.0, 4.1 VMware ESXi 4.0, 4.1 VMware ESXi 5.0, 5.1 Windows 2008 R2 with Hyper-V

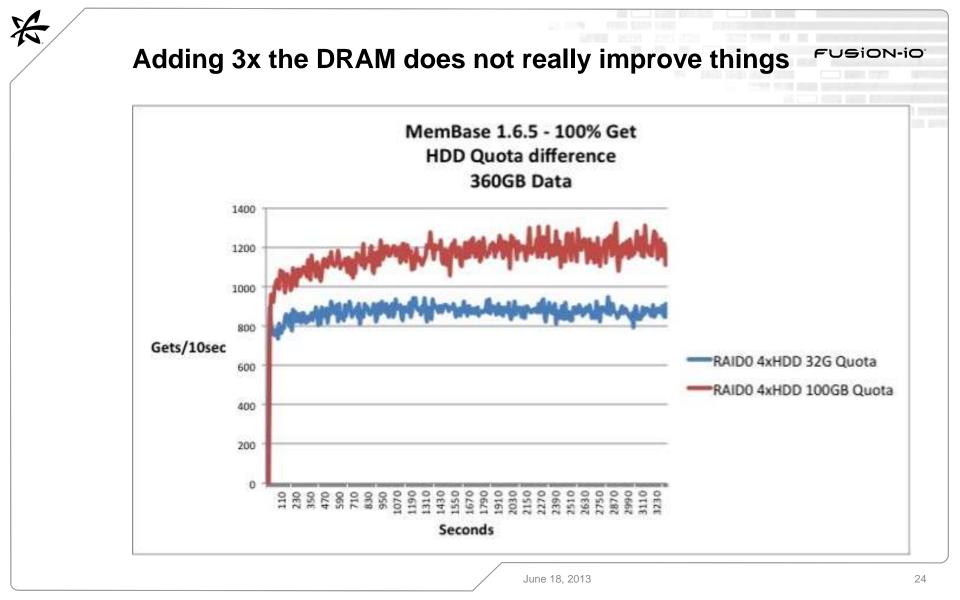
Unix

Solaris 10 x64 U8, U9, U10 OpenSolaris 2009.06 x64 OSX 10.6 and later FreeBSD 8,9











HBase Server

A typical server...

CPU Cores: 32 with HT Memory: 128 GB



Is your working set larger than 128GB?



With NoSQL Databases, we tend to scale out for DRAM

Combined Resources CPU Cores: 96 Memory: 384 GB



More cores than needed to serve reads and writes.

The HBase BucketCache (HBase-7404)

Committed to HBase trunk. Will be in 0.96 release, backport patch for 0.94 available.`

Victim cache for LRUBlockCache – Move fast ioMemory close to DRAM cache



https://issues.apache.org/jira/browse/HBASE-7404



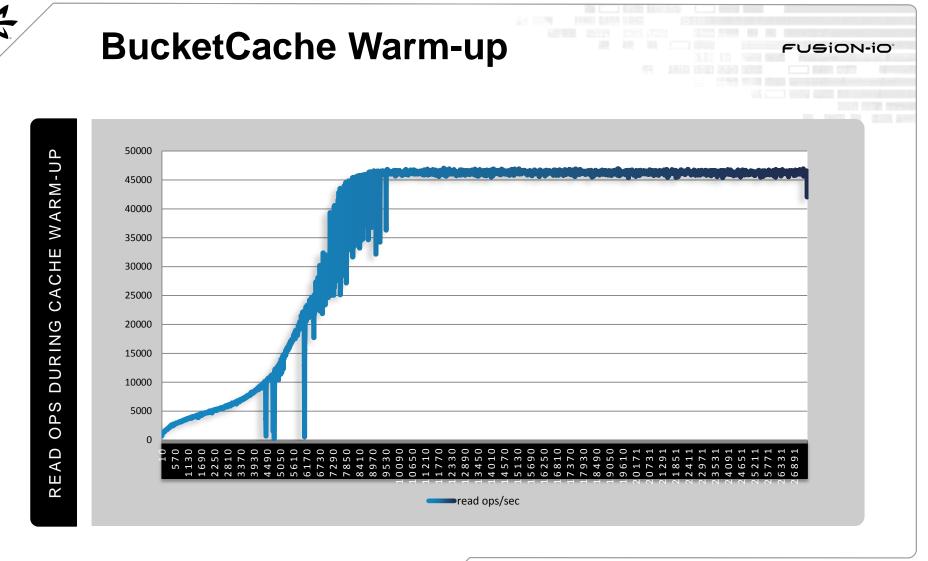
BucketCache Configuration

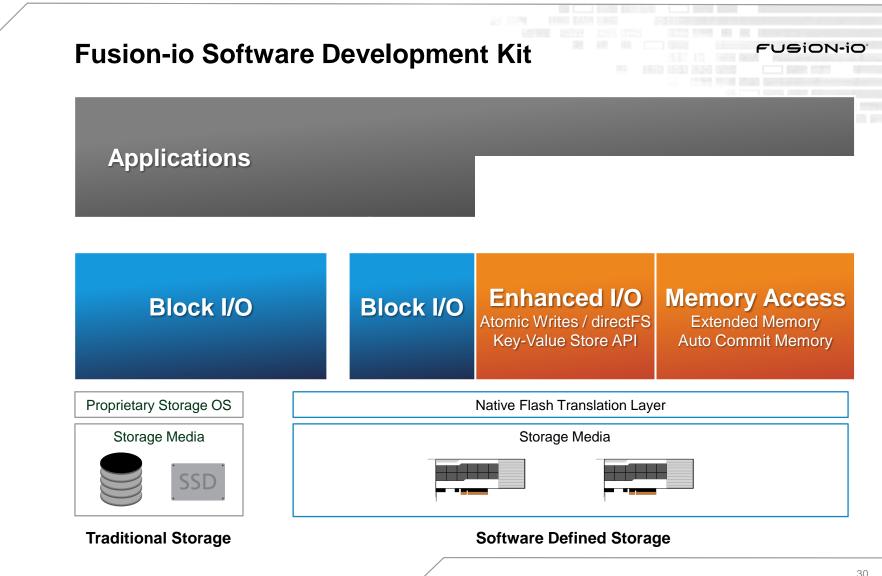
In hbase-site.xml

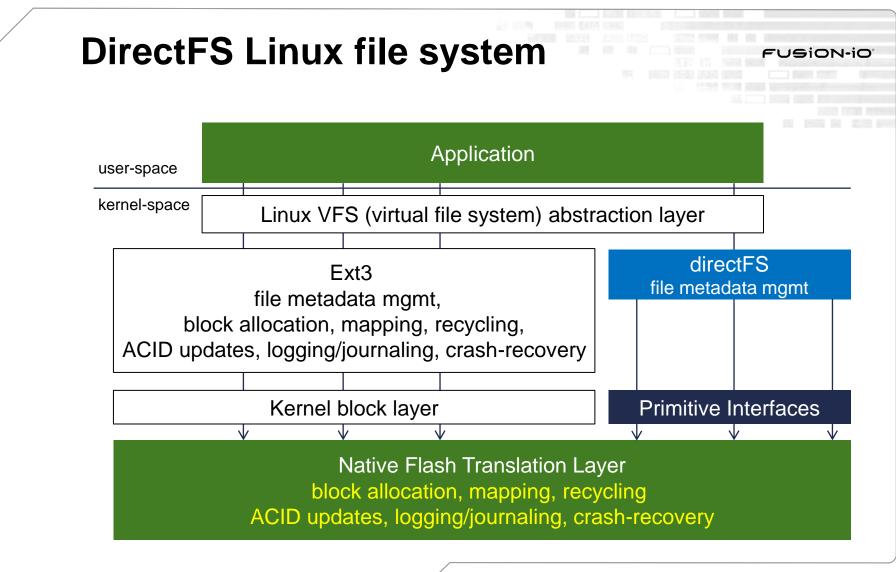
<property> <name>hbase.bucketcache.ioengine</name> <value>file:/path /to/bucketcache.dat</value> </property> <property> <name>hbase.bucketcache.size</name> <!-- 2TB: unit is MB -->

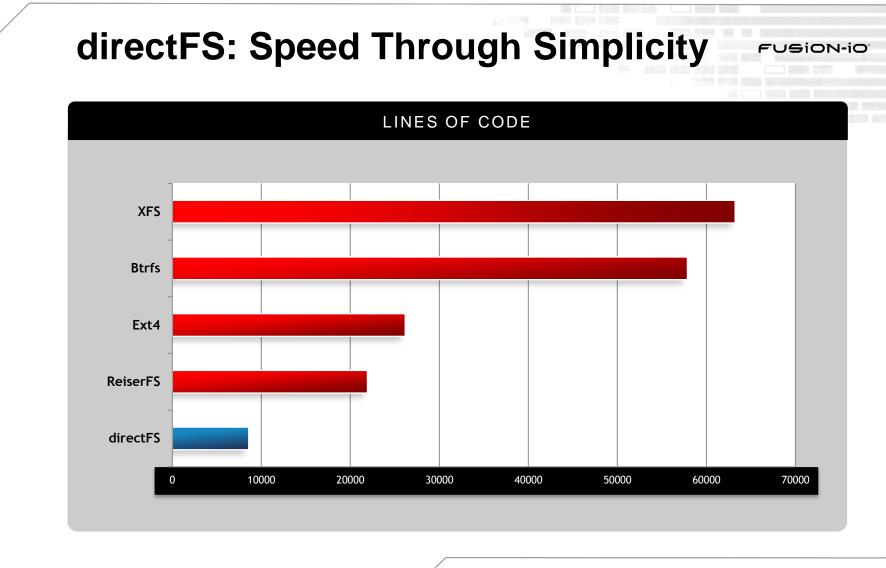
- <value>2097152</value>
- </property>

FUSION-iO









Atomic writes – Transactional I/O

- System call tells DirectFS that all I/O to this file should be treated as atomic
- Avoids the partial page write problem
- Accepted by T10 technical committee for SCSI standard
- Minimal application changes required

Percona Server, MariaDB, MySQL 5.6 FUSION-IO

- Efficient XtraDB/InnoDB storage engine
- Well optimized for seek-less storage like flash
- Many config parameters to fine-tune performance

- What else can be done?
 - Lock contention can still be improved as seen by using multiple instances with the same storage device
 - Tapping into the native performance of flash by exposing key FTL features to the application

MySQL Writes Comparison

Traditional MySQL Writes

MySQL with Atomic Writes

Page

Page Page

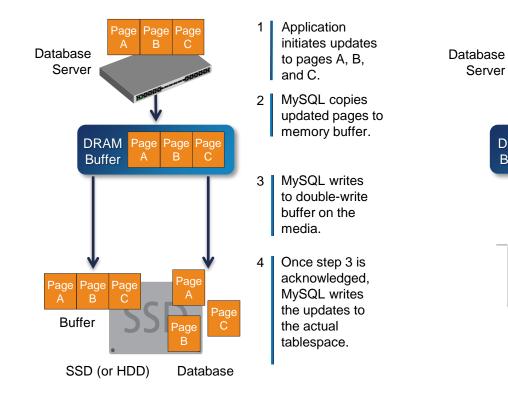
Page

Page

DRAM Page

ioMemory

Buffer



Application initiates updates to pages A, B, and C.

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MySQL copies updated pages to memory buffer.

2

3

Page

Database

MySQL writes to actual tablespace, bypassing the double-write buffer step due to inherent atomicity guaranteed by the intelligent device.



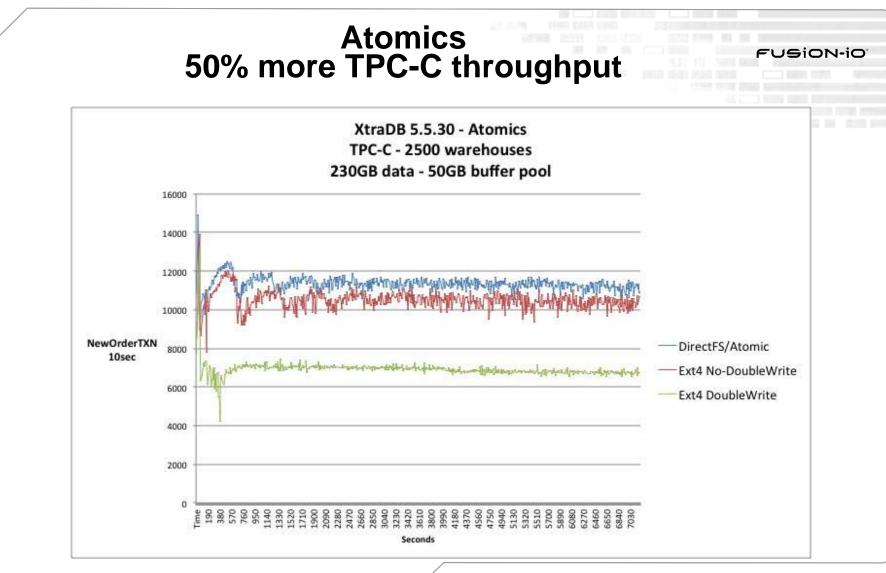
Atomic benchmarks

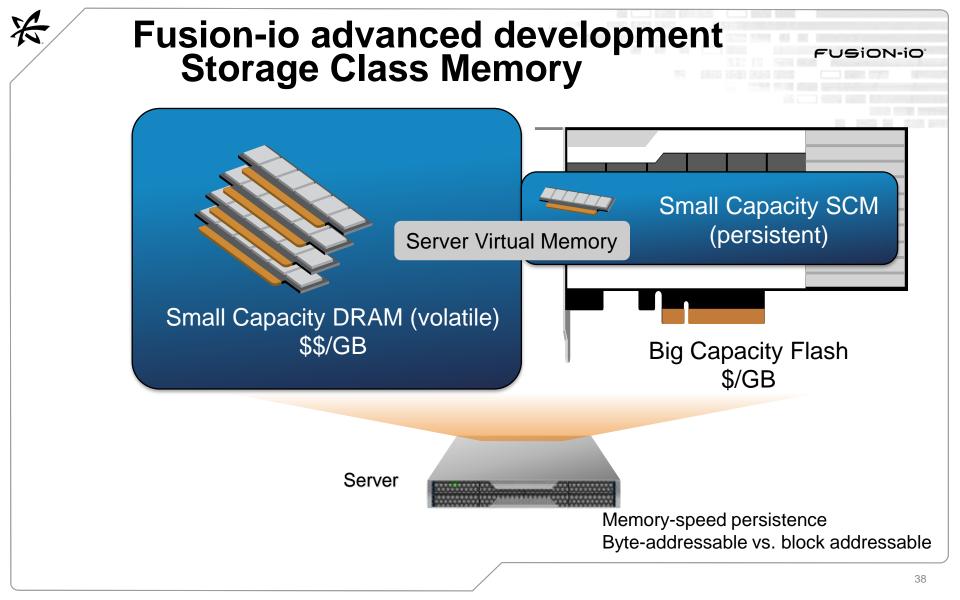
First, lets sum up the MySQL benefits here:

 Writing only 50% of the data otherwise required for ACID compliance

That's pretty much it...but it gives us

- Twice the flash endurance
- Much better latency because of fewer syscalls
- Much better application throughput due to less I/O
- Better concurrency due to fewer locks







 Lets look at keeping a database log using memory semantics

Goal is to reduce latency, cost of flushing data to a persistent state and further minimize writes

SCM testing using modified Innosim toool

SCM logger interface

> logger_open()

Open and initialize logging infrastructure within the FTL

logger_close()

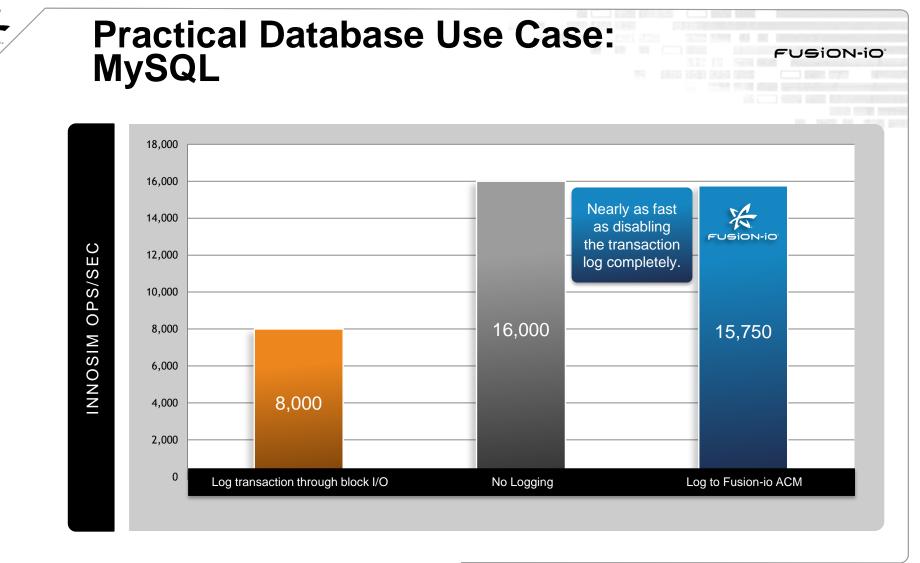
Clean-up

logger_append()

Append to head of log at memory speeds. This basically translates to a memcpy()

Iogger_sync()

Serialize data using assembler 'mfence' instruction





As an SSD, flash accelerates applications.

At full maturity, Non-Volatile Memory will transform software development.

Native Flash API availability

- Percona Server: 5.5.31
- MariaDB mainline: 5.5.31
- Oracle MySQL:
 - https://code.launchpad.net/~tmathiasen/mysql-server/mysql-5.5-fio
- Cassandra atomics implementation in progress
- DirectFS public beta expected July 22nd

THANK YOU

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