

Dynamo concepts in depth.
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The shopping cart case



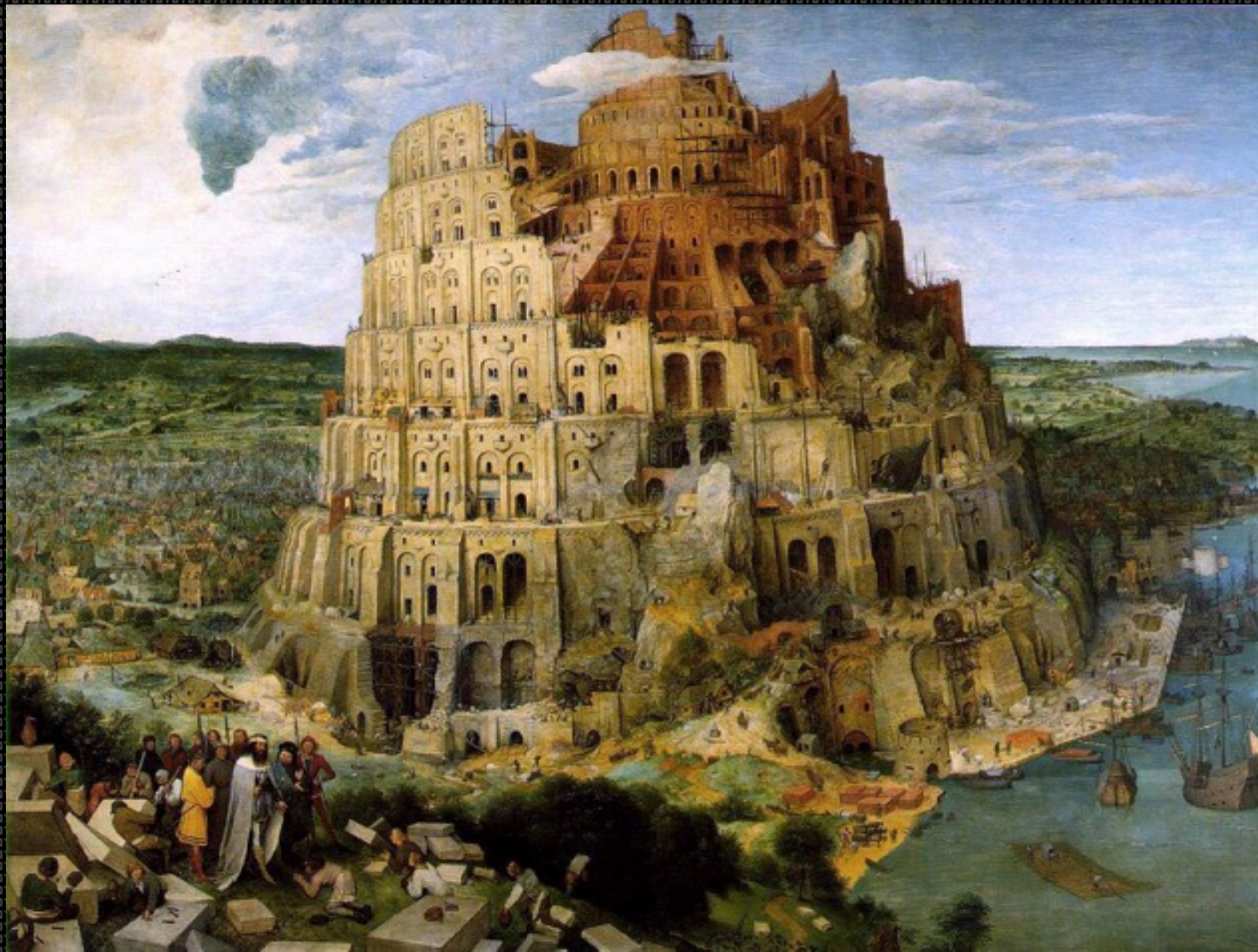
Friday, August 31, 12

The 2 AM alarm call case



Friday, August 31, 12

The Tower of Babel case



The Neo vs. Smiths case



The Pavlo case



Friday, August 31, 12



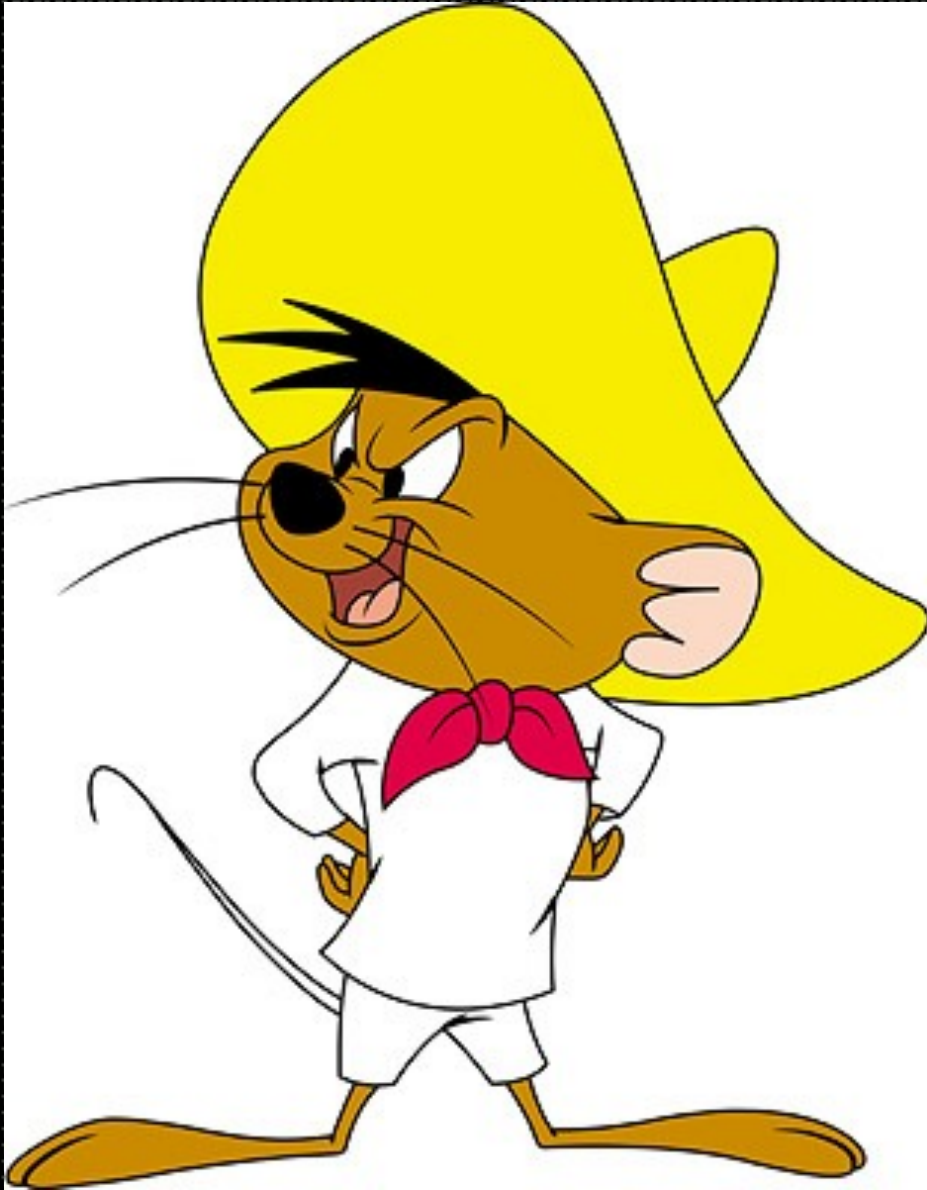
So Dynamo isn't about speed.

It's about immediate,
reliable writes.

It's about
operation relaxation.

It's about distribution
and fault tolerance.

It's about almost
linear scalability.



Time and timestamps

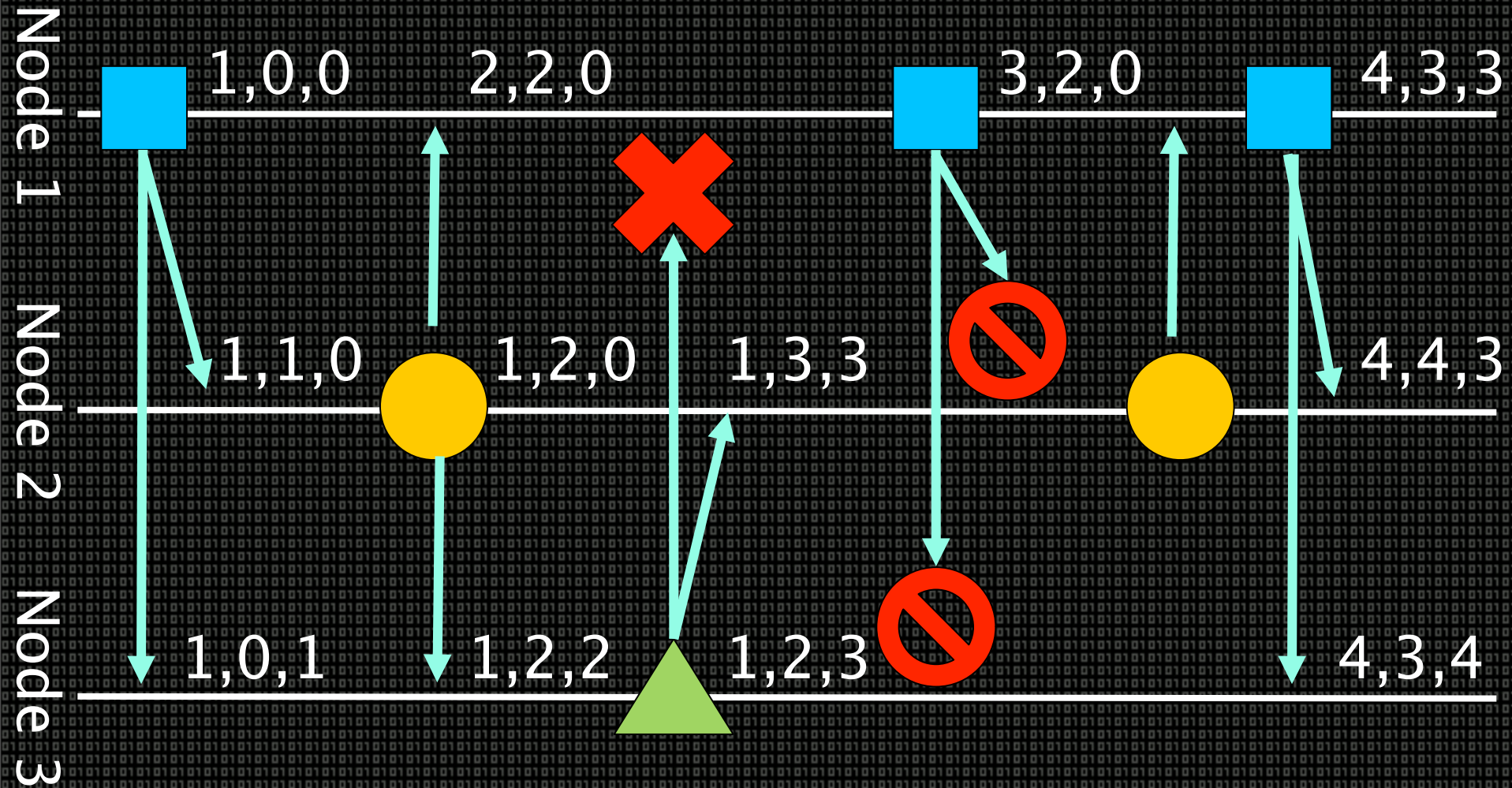


$V(i)$, $V(j)$: competing

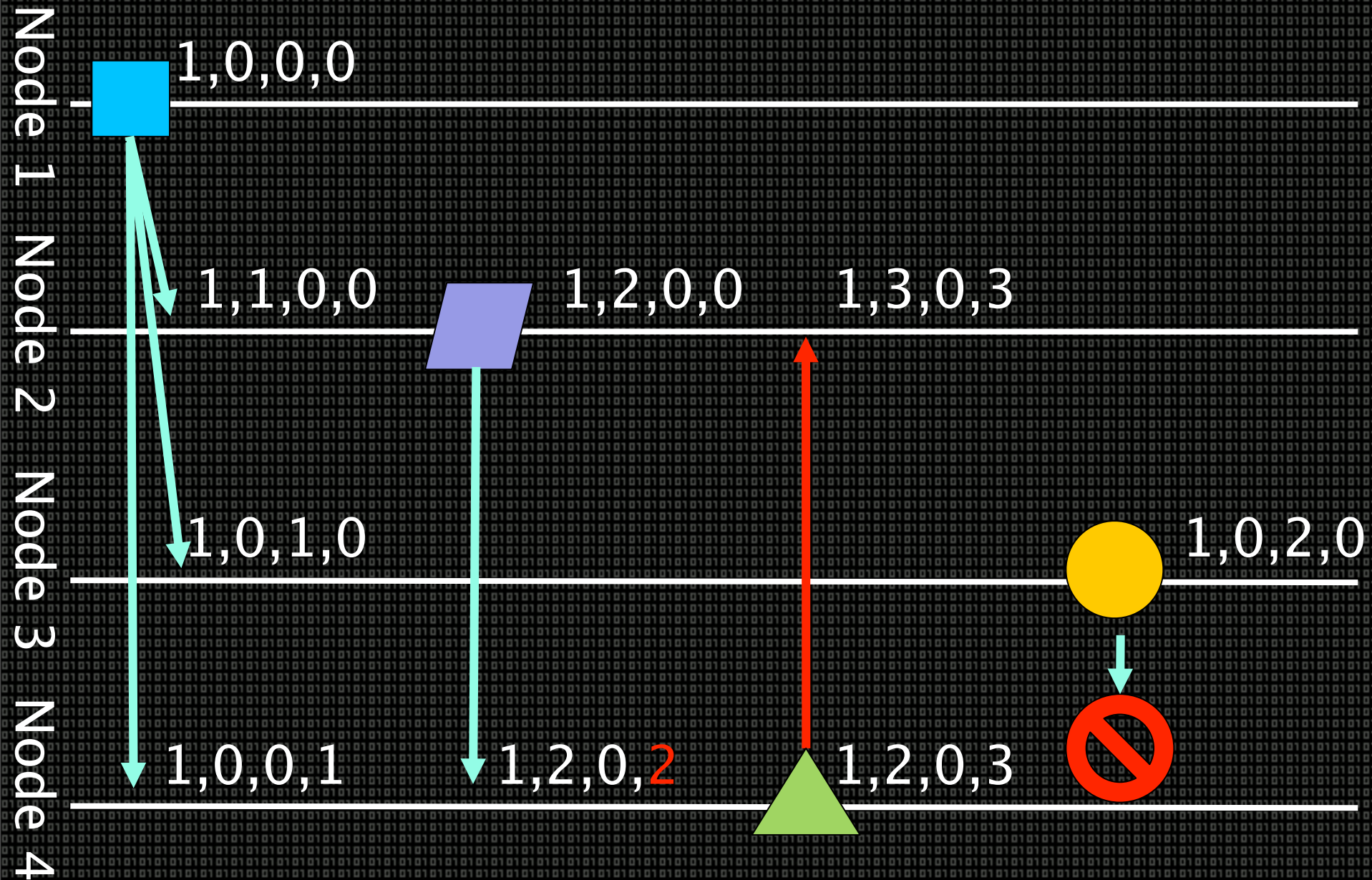
Conflict resolution:

- 1: **siblings**, client
- 2: **merge**, system
- 3: **voting**, system

Vector clocks



Vector clocks



$O(1)$ for data lookups / delta tracking

#

Merkle Trees

N, M: nodes

HT(N), HT(M): hash trees

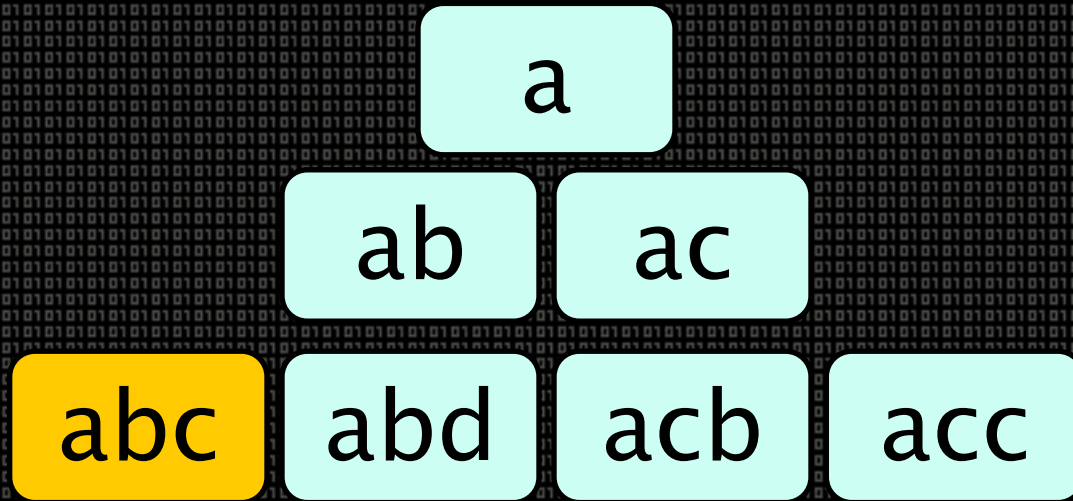
M needs update:

- obtain HT(N)

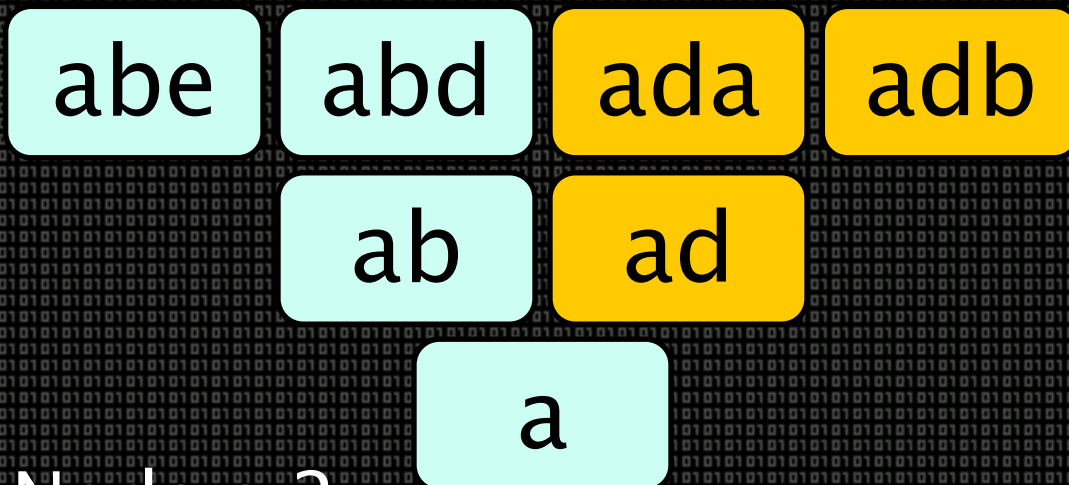
- calc $\text{delta}(\text{HT}(\text{M}), \text{HT}(\text{N}))$

- pull keys(delta)

Node a.1



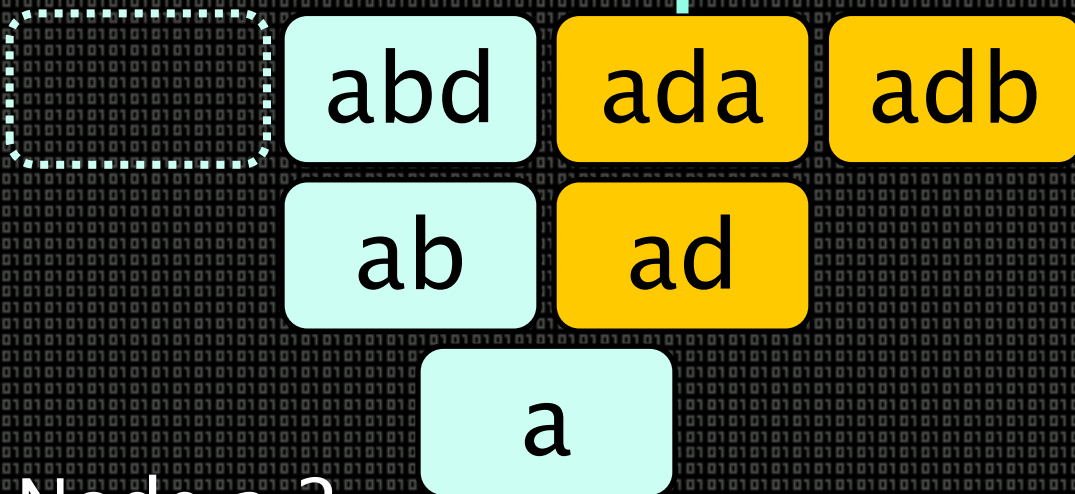
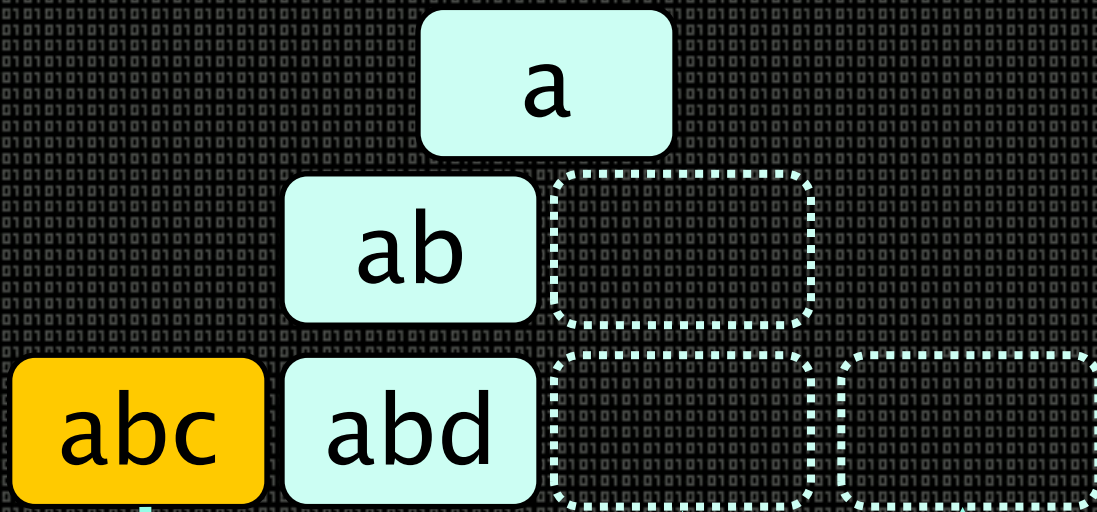
Merkle Trees



Node a.2

Node a.1

Merkle Trees



Node a.2

"Equal" nodes based decentralized distribution



Consensus, agreement, voting, quorum



Consistent hashing – the ring

X bit integer space

$$0 \leq N \leq 2^X$$

or: $2 \times \pi$

$$0 \leq A \leq 2 \times \pi$$

$$x(N) = \cos(A)$$

$$y(N) = \sin(A)$$

Quorum

V: vnodes holding a key

W: write quorum

R: read quorum

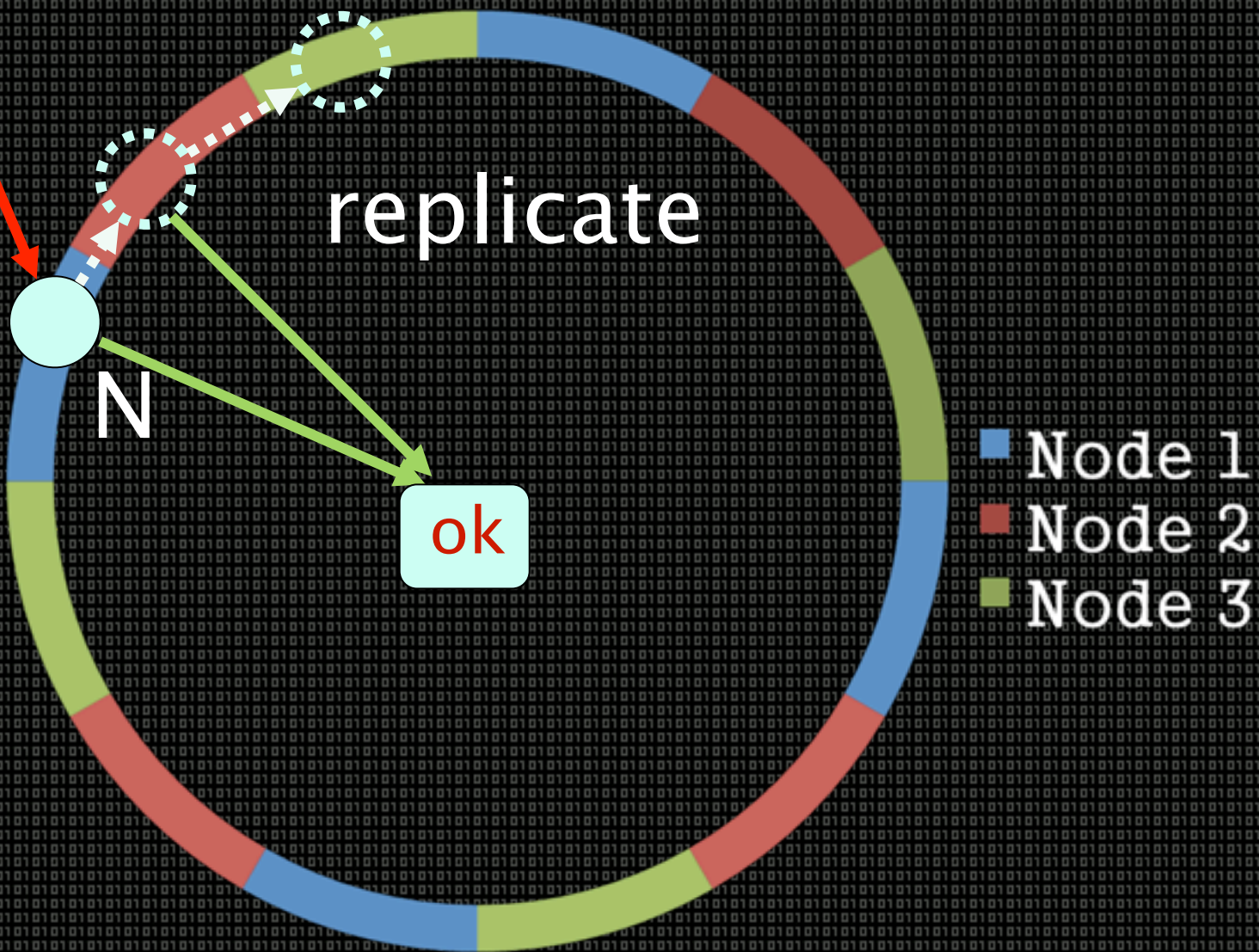
DW: durable write quorum

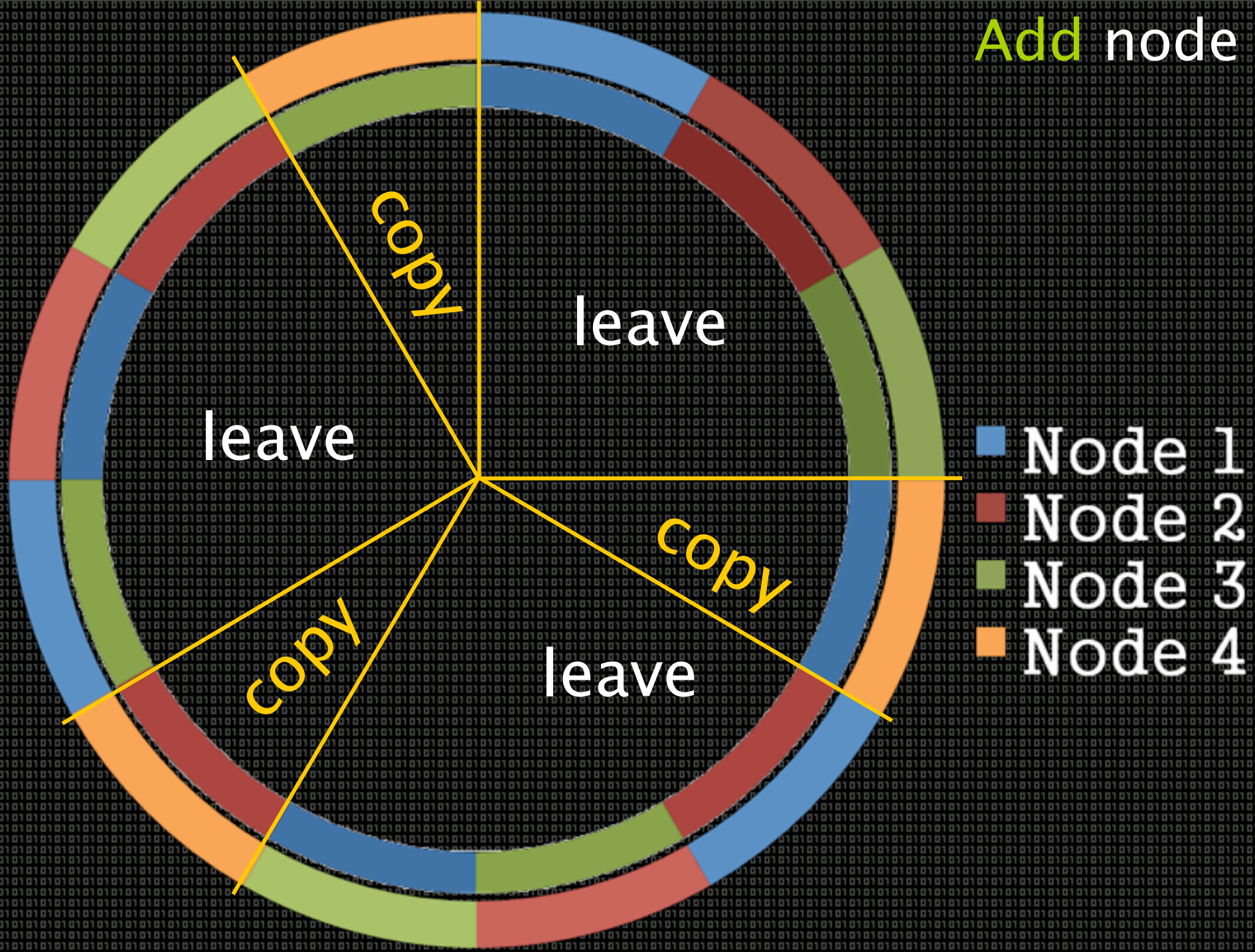
$$W > 0.5 * V$$

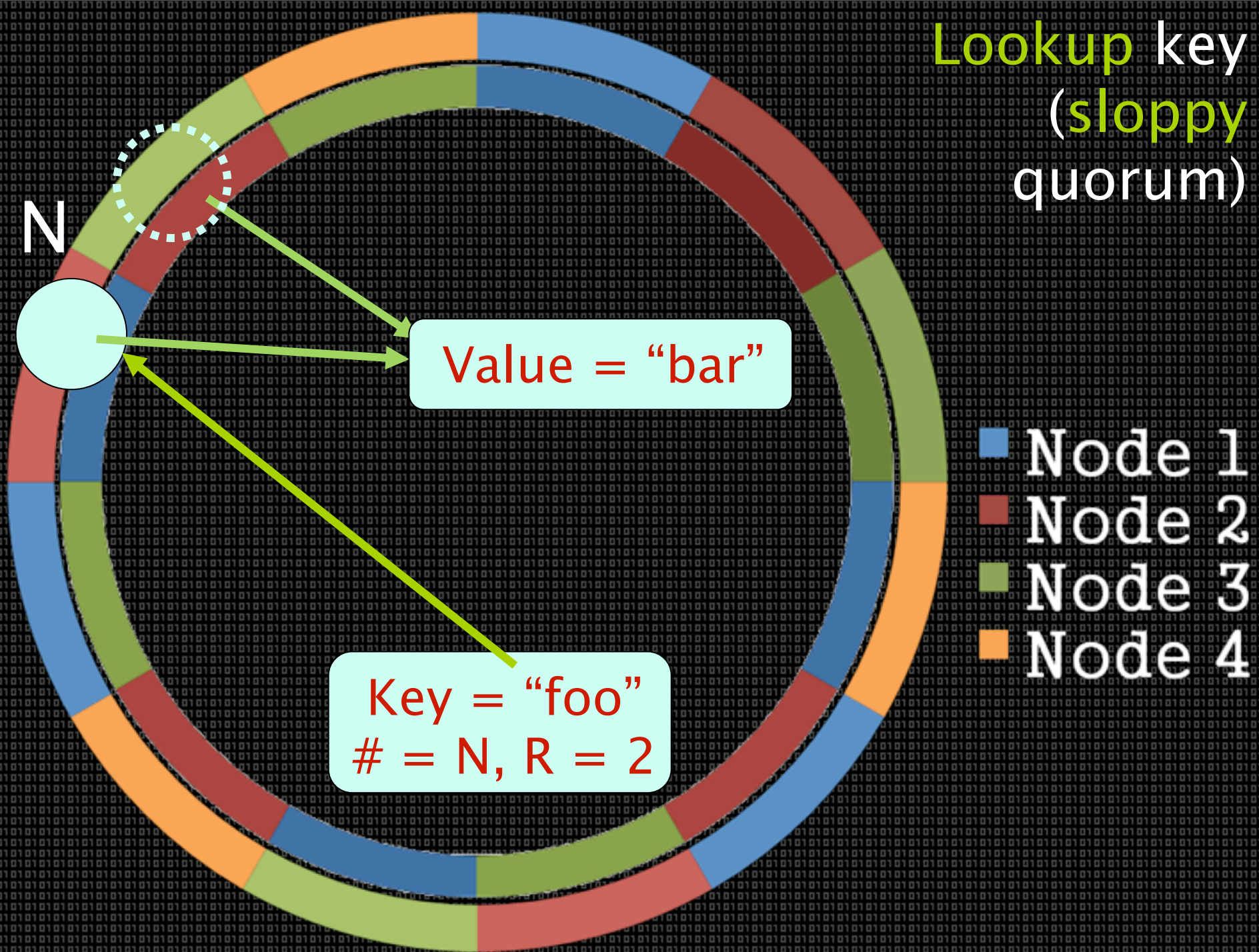
$$R + W > V$$

Key = "foo"
= N, W = 2

Insert key
(sloppy quorum)







Remove
node

copy

leave

■ Node 1
■ Node 2
■ Node 3

Gossip – node down/up

Node 1
Node 2
Node 3
Node 4

update

update,
4 down

update

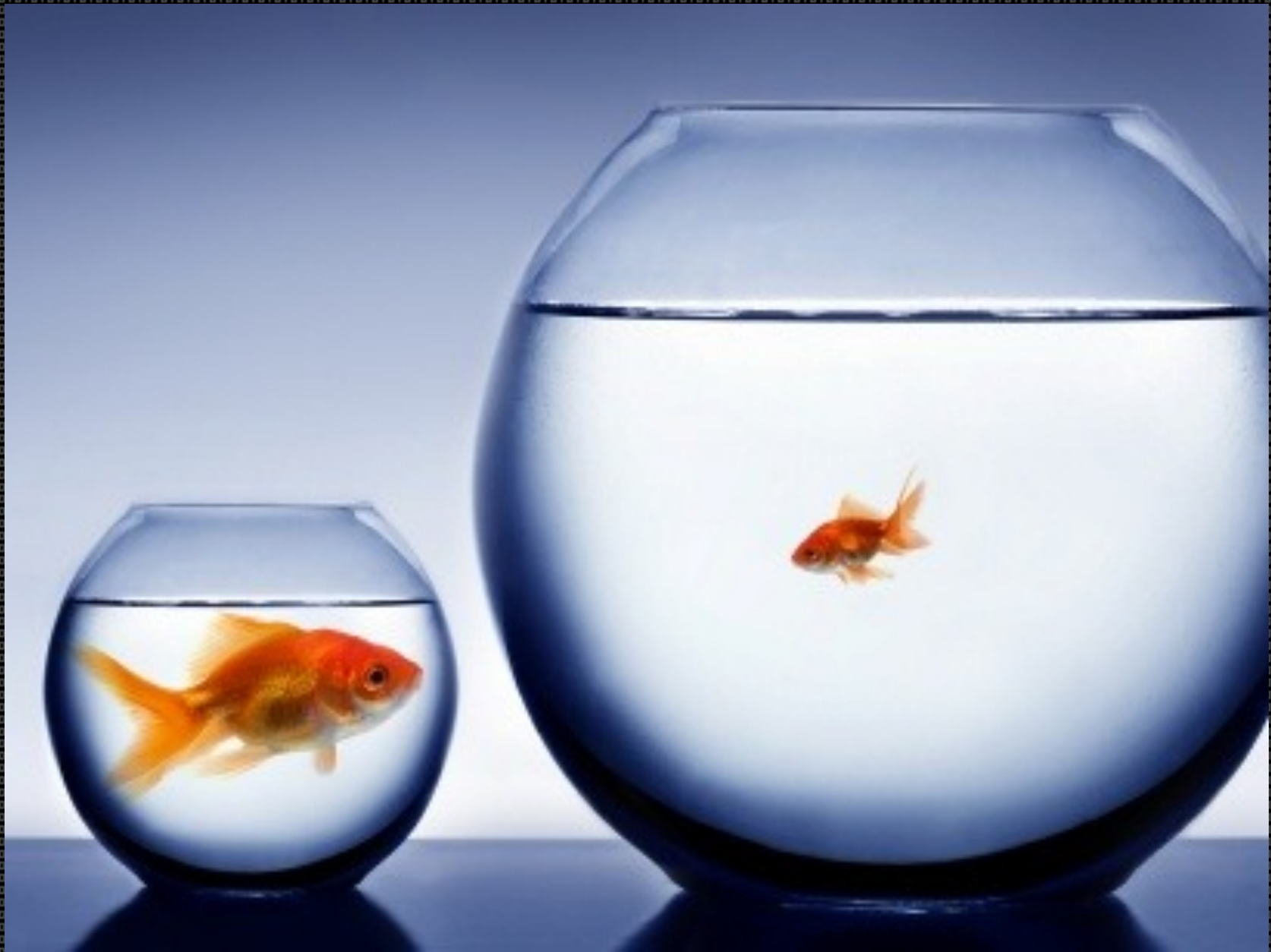
read,
4 up

update

read



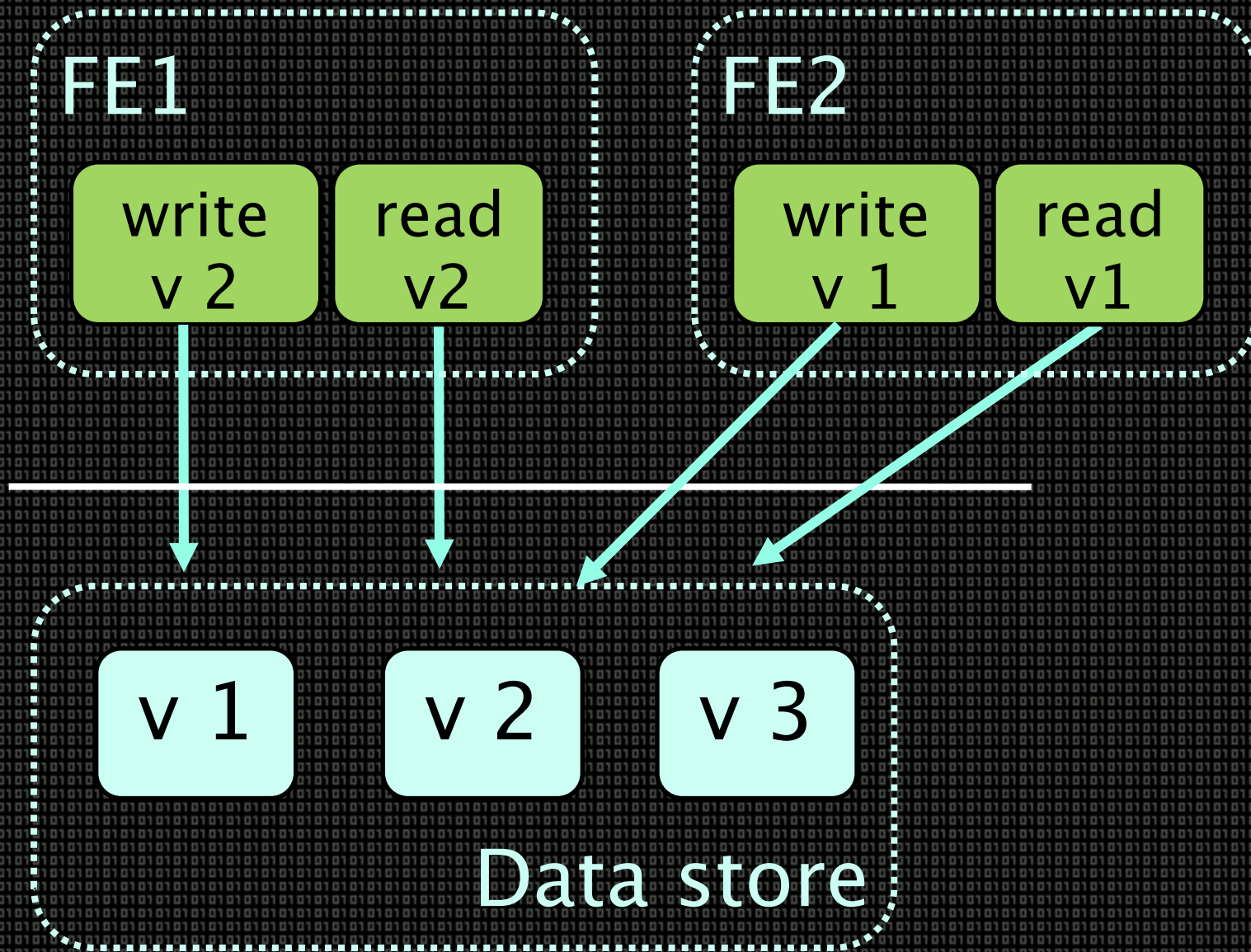
Eventual consistency



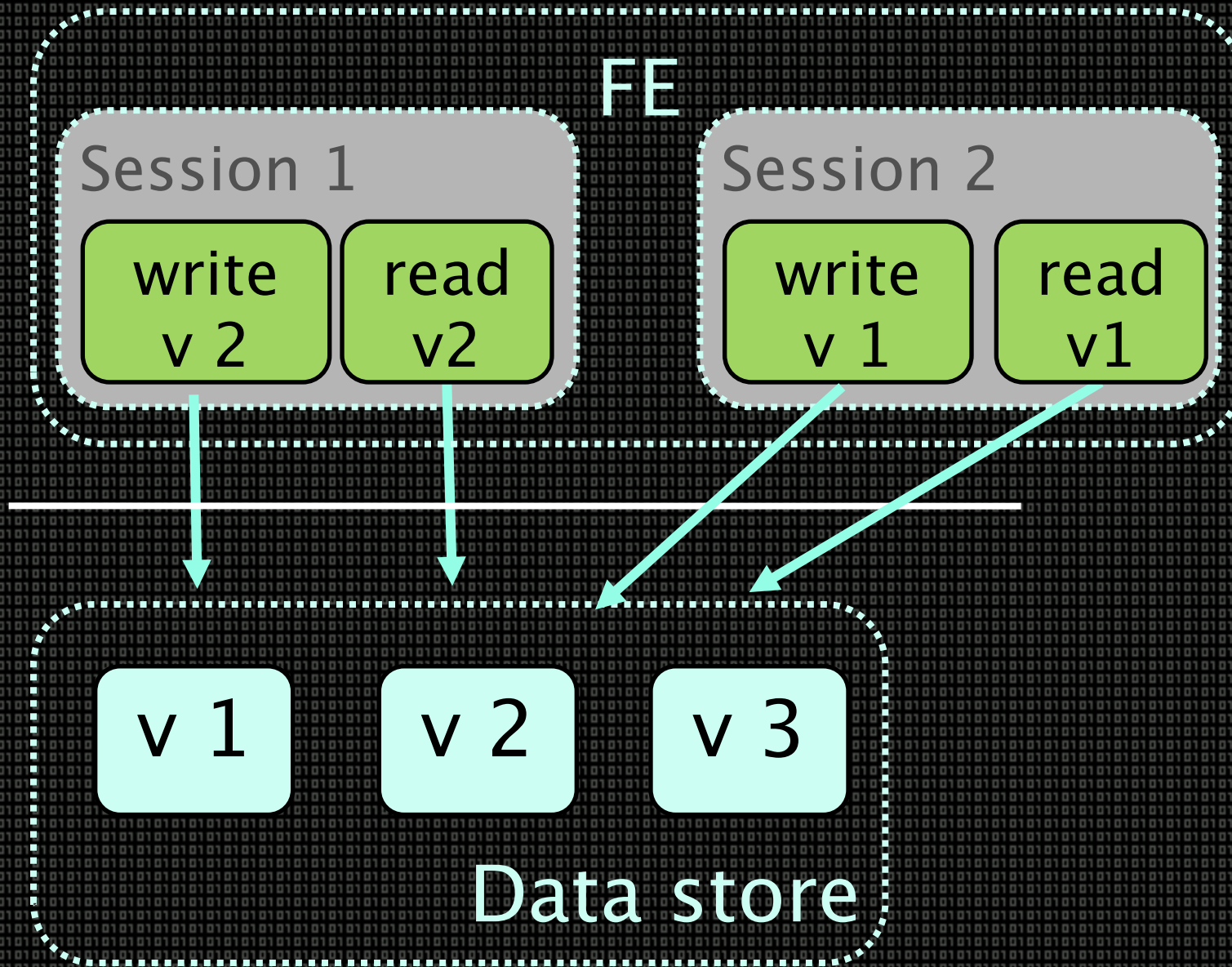
Basically Available,
Soft-state,
Eventually consistent

Opposite to ACID

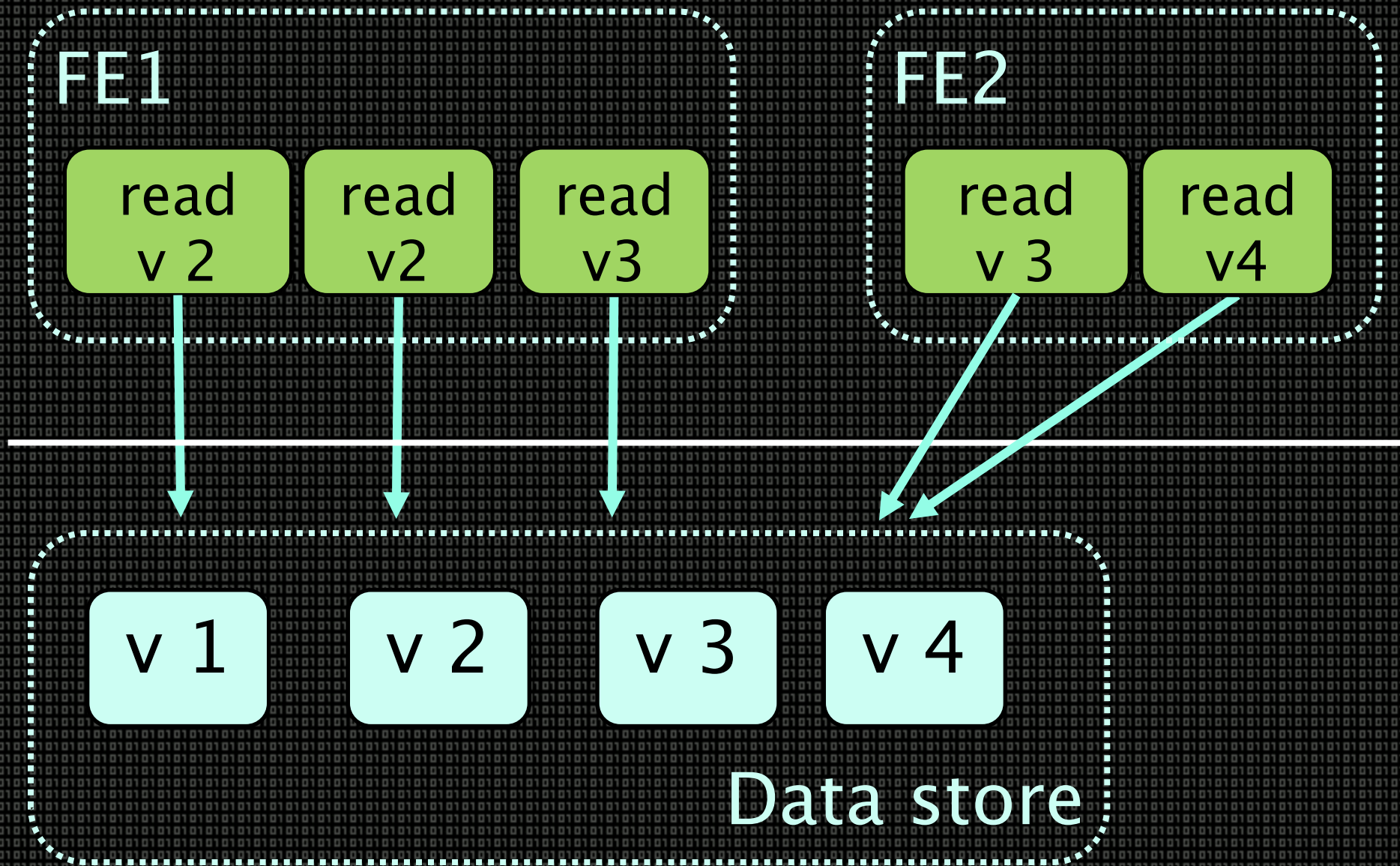
Read **your** write consistency



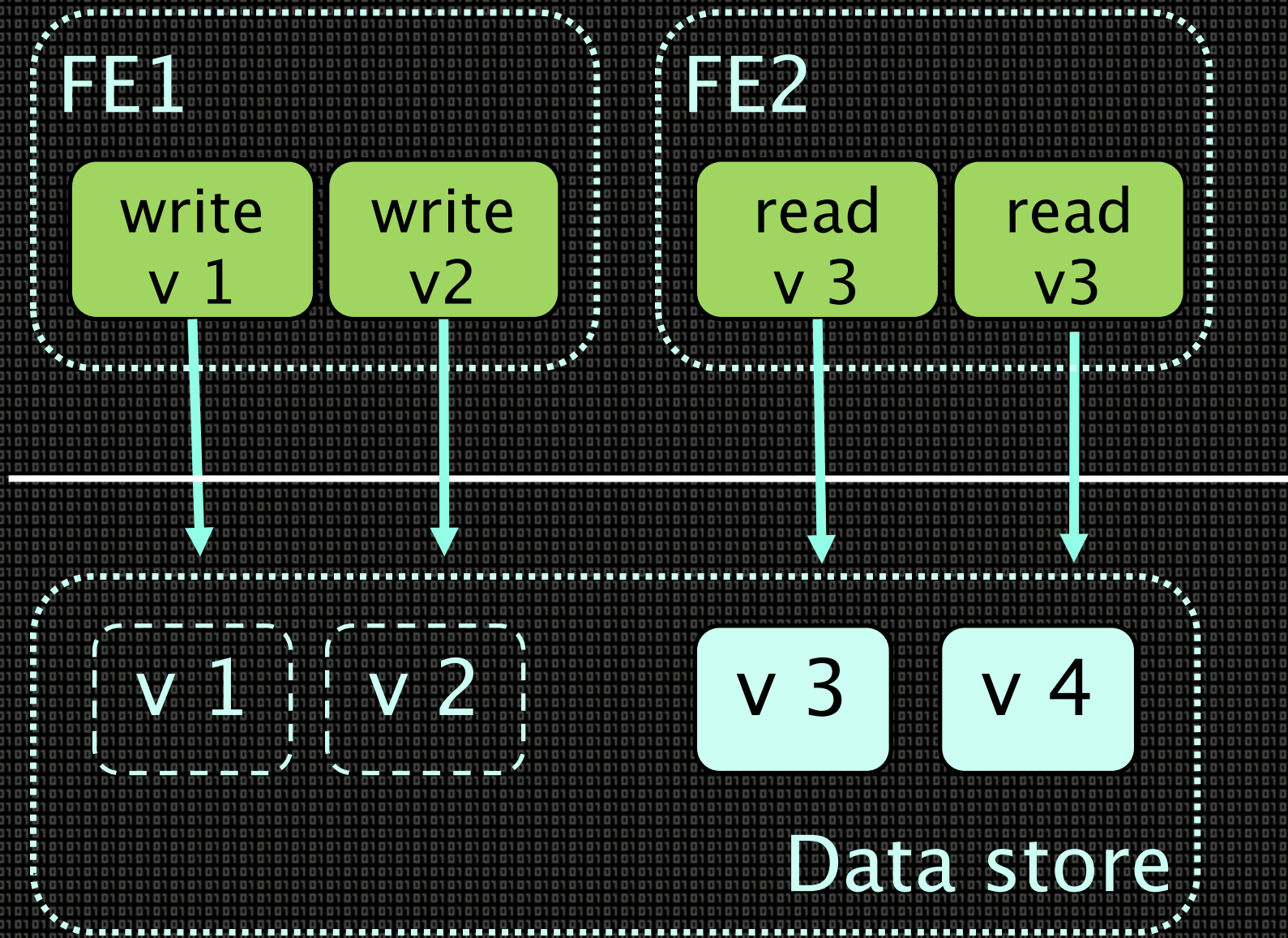
Session consistency



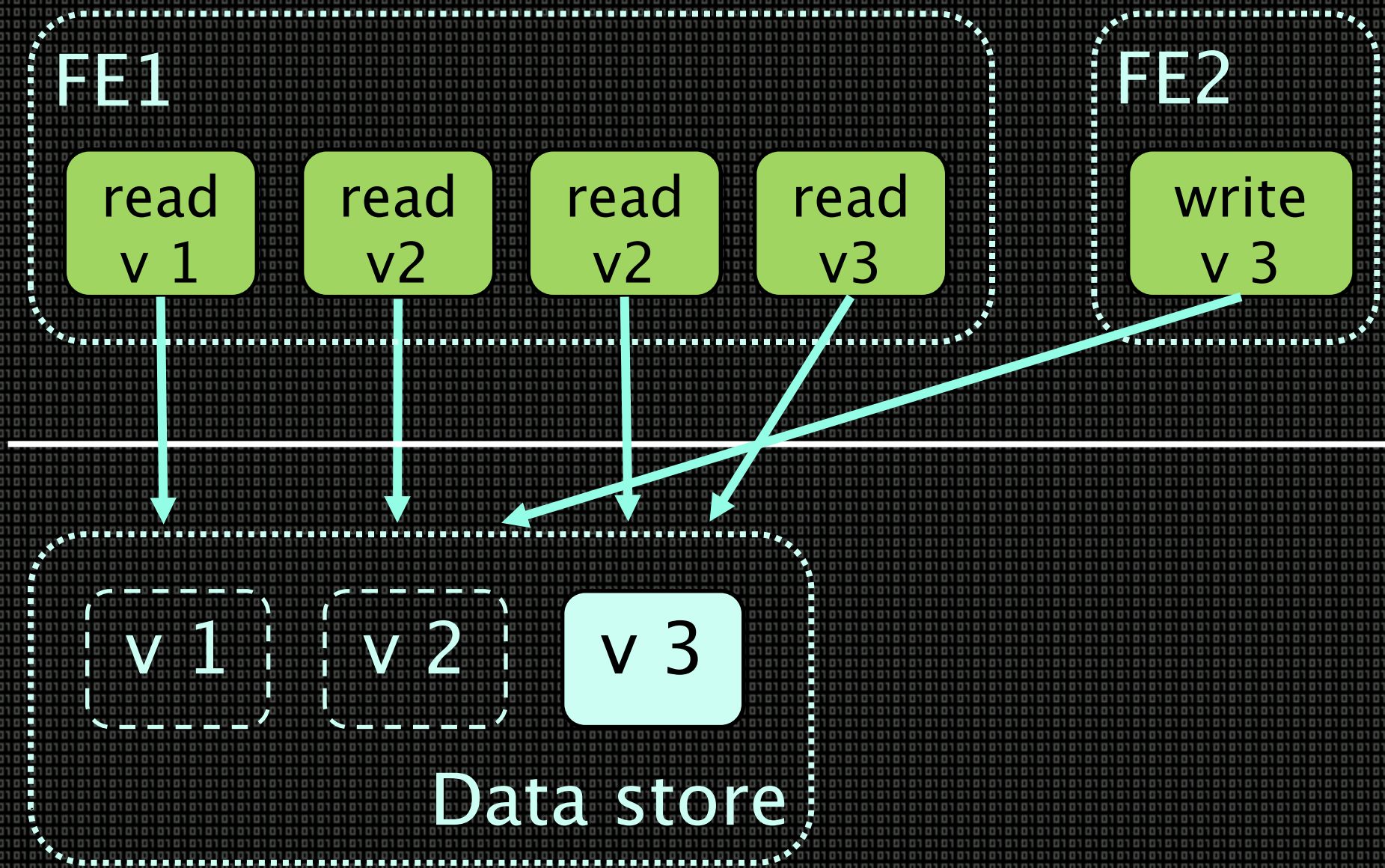
Monotonic **read** consistency



Monotonic **write** consistency



Eventual consistency



N: node, G: group including N

node(N) is unavailable

replicate to G or

store data(N) locally

hint handoff for later

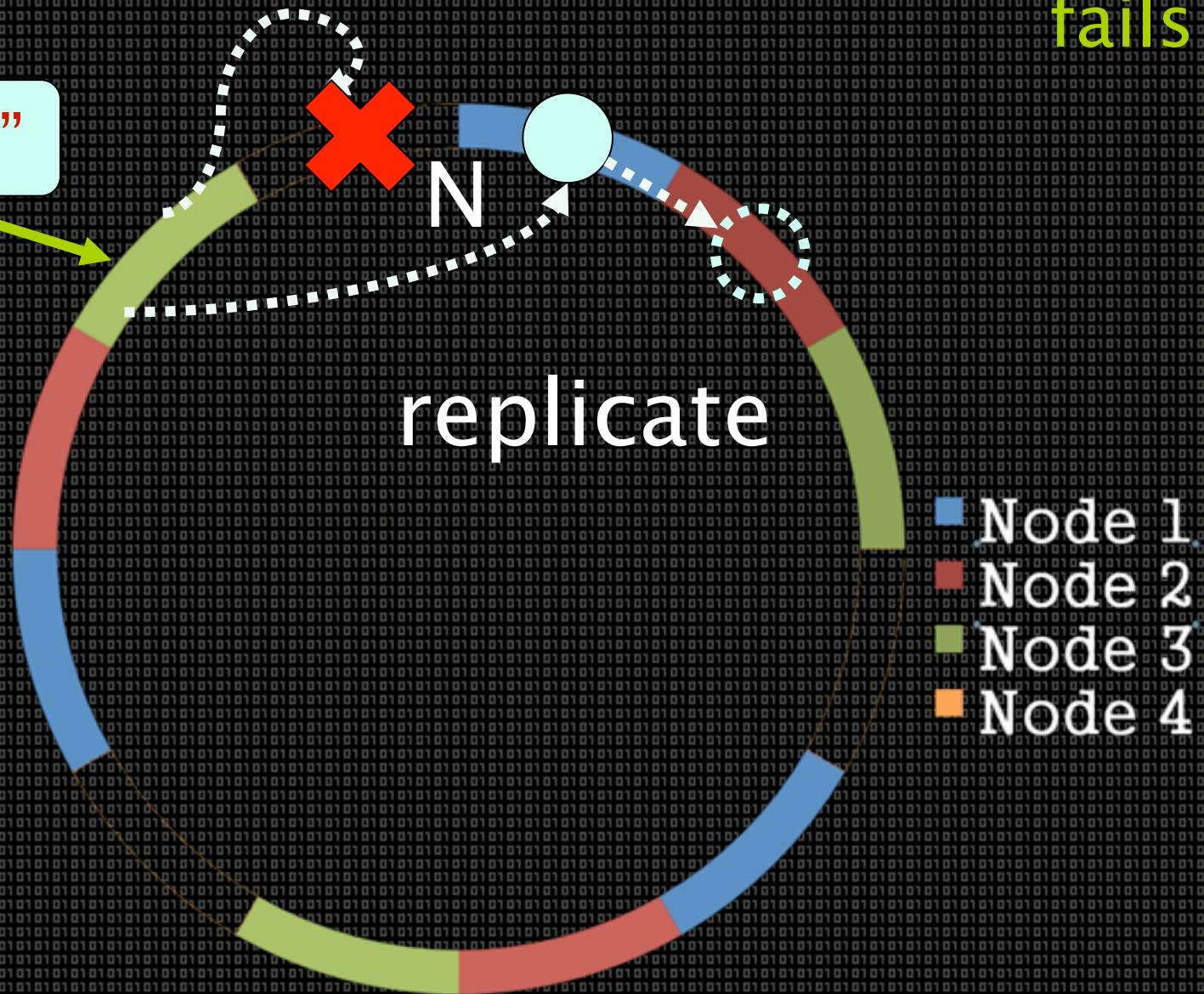
node(N) is alive

handoff data to node(N)

Key = "foo", # = N ->
handoff hint = true

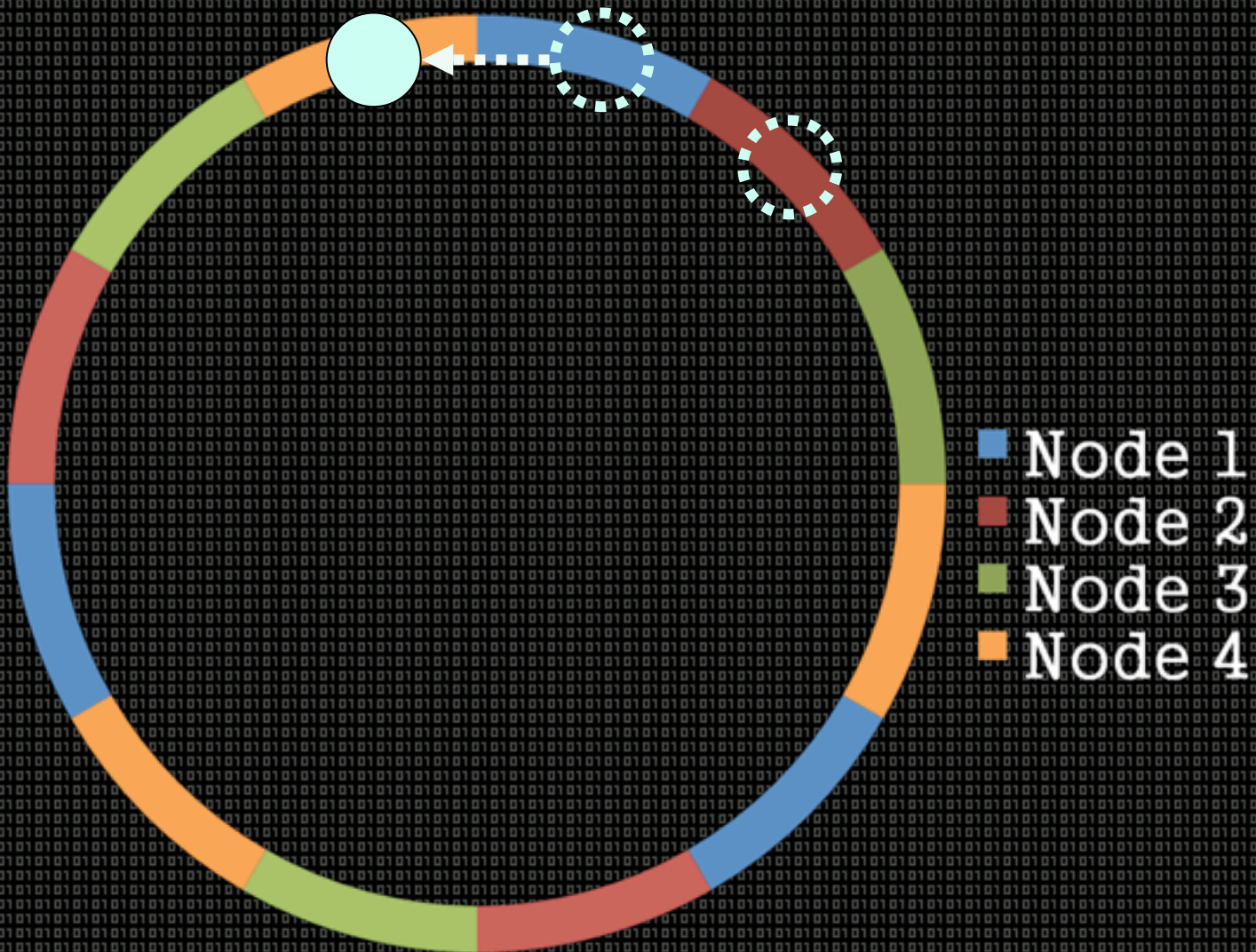
Direct
replica
fails

Key = "foo"



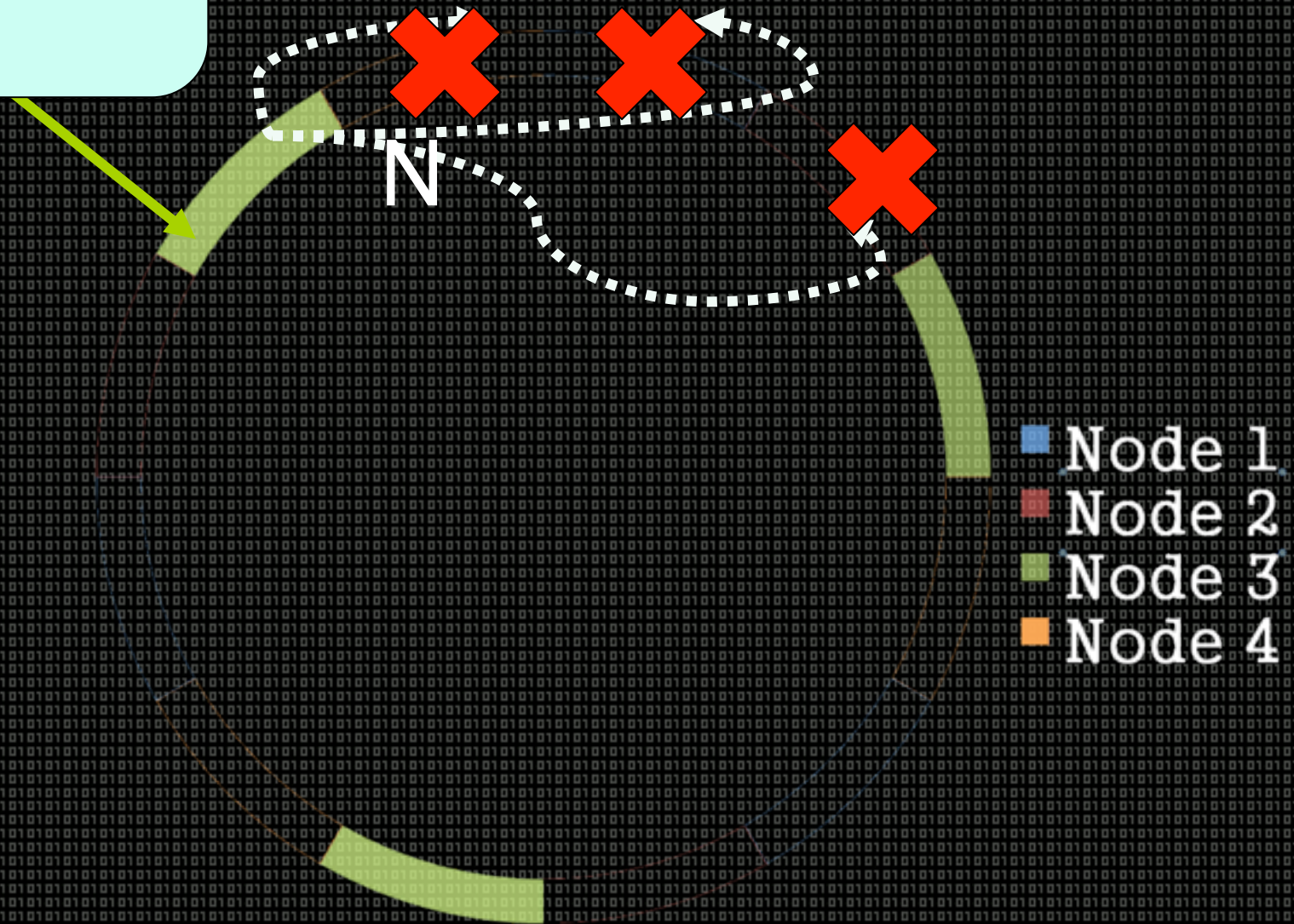
handoff

Replica
recovers



Key = "foo",
= N ->
handoff hint =
true

All
replicas
fail

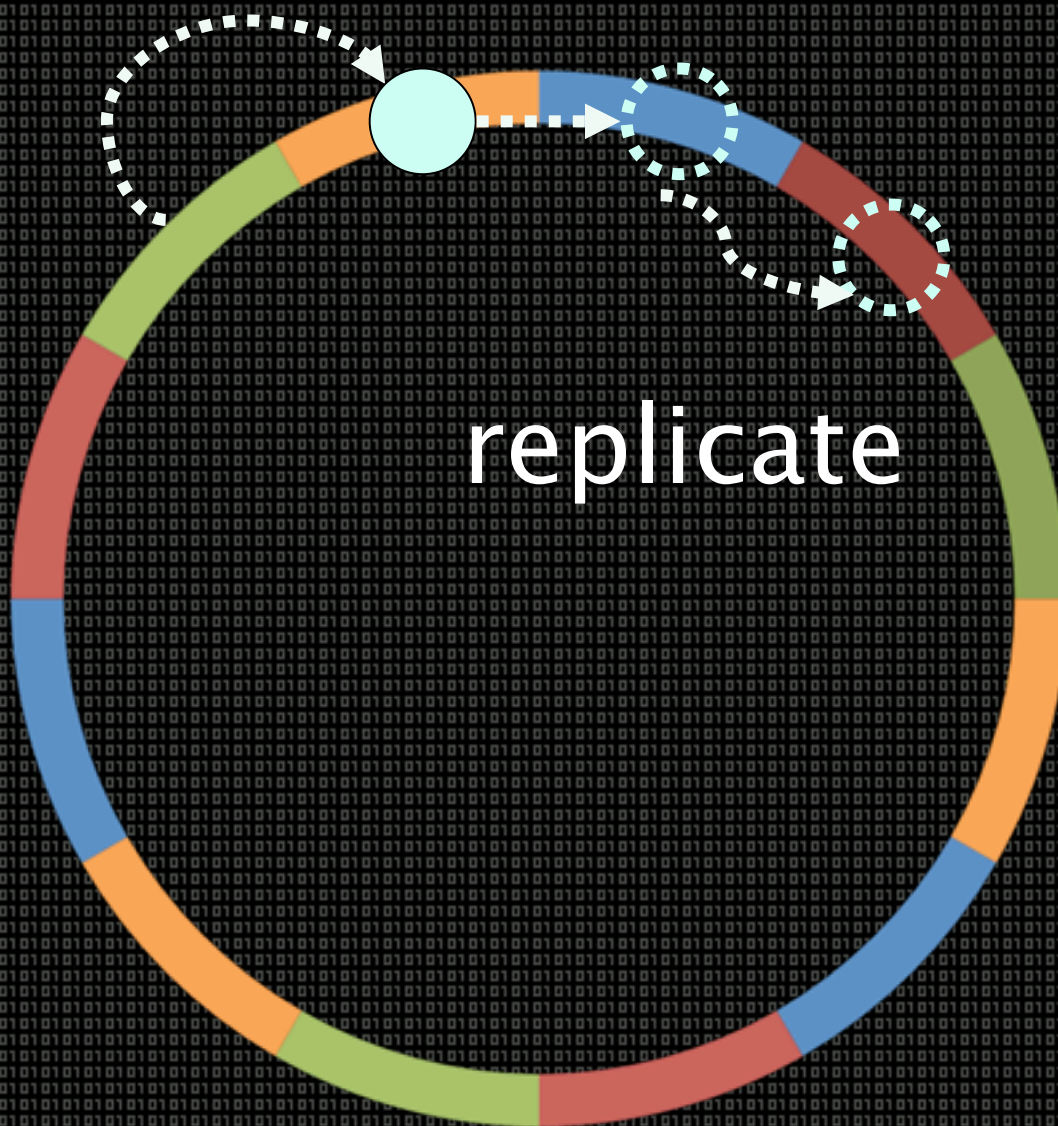


All

replicas

recover

handoff



- Node 1
- Node 2
- Node 3
- Node 4



Latency is an adjustment screw



Availability is an adjustment screw



CAP – the variations

CA – irrelevant

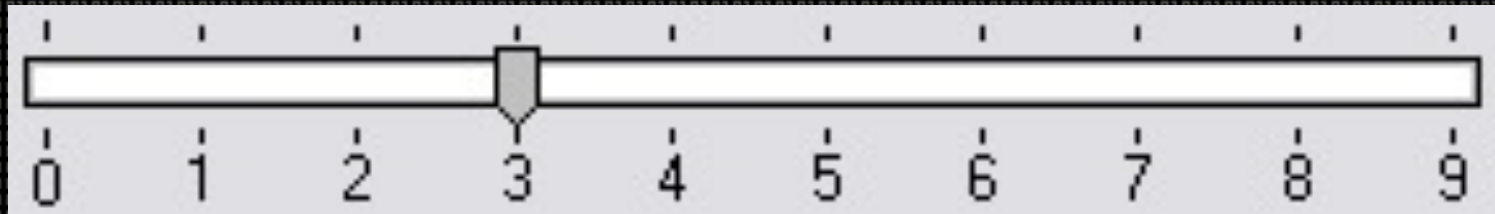
CP – eventually unavailable offering maximum consistency

AP – eventually inconsistent offering maximum availability

CAP – the tradeoff

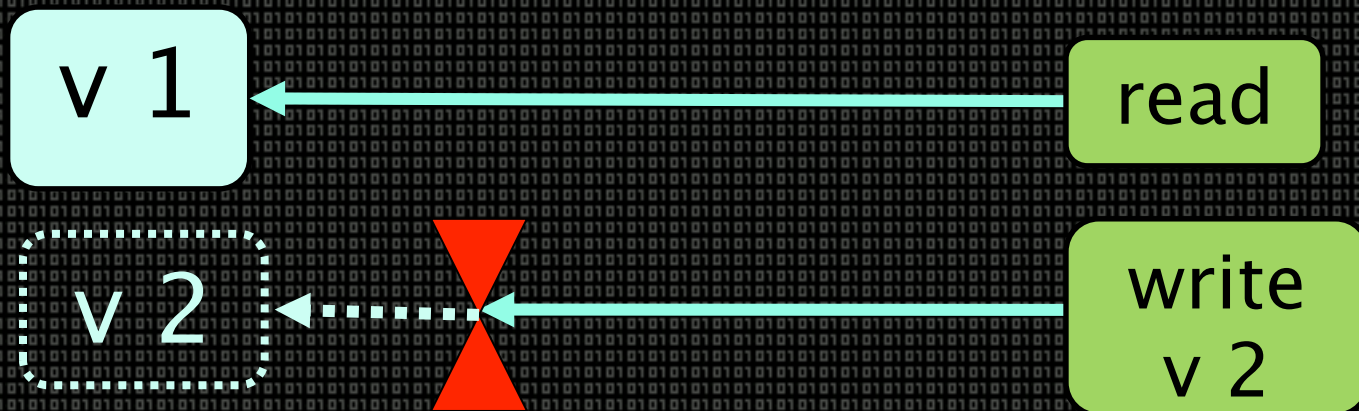
A

C

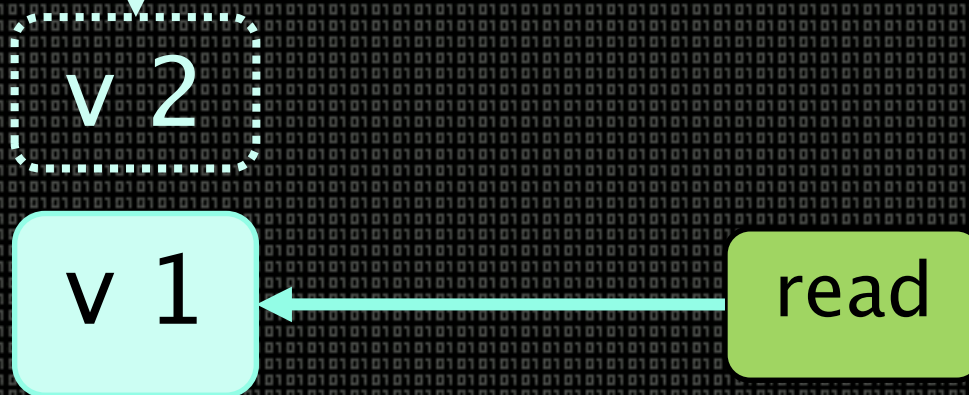


Replica 1

CP



Replica 2



Replica 1

CP (partition)

v 1

read

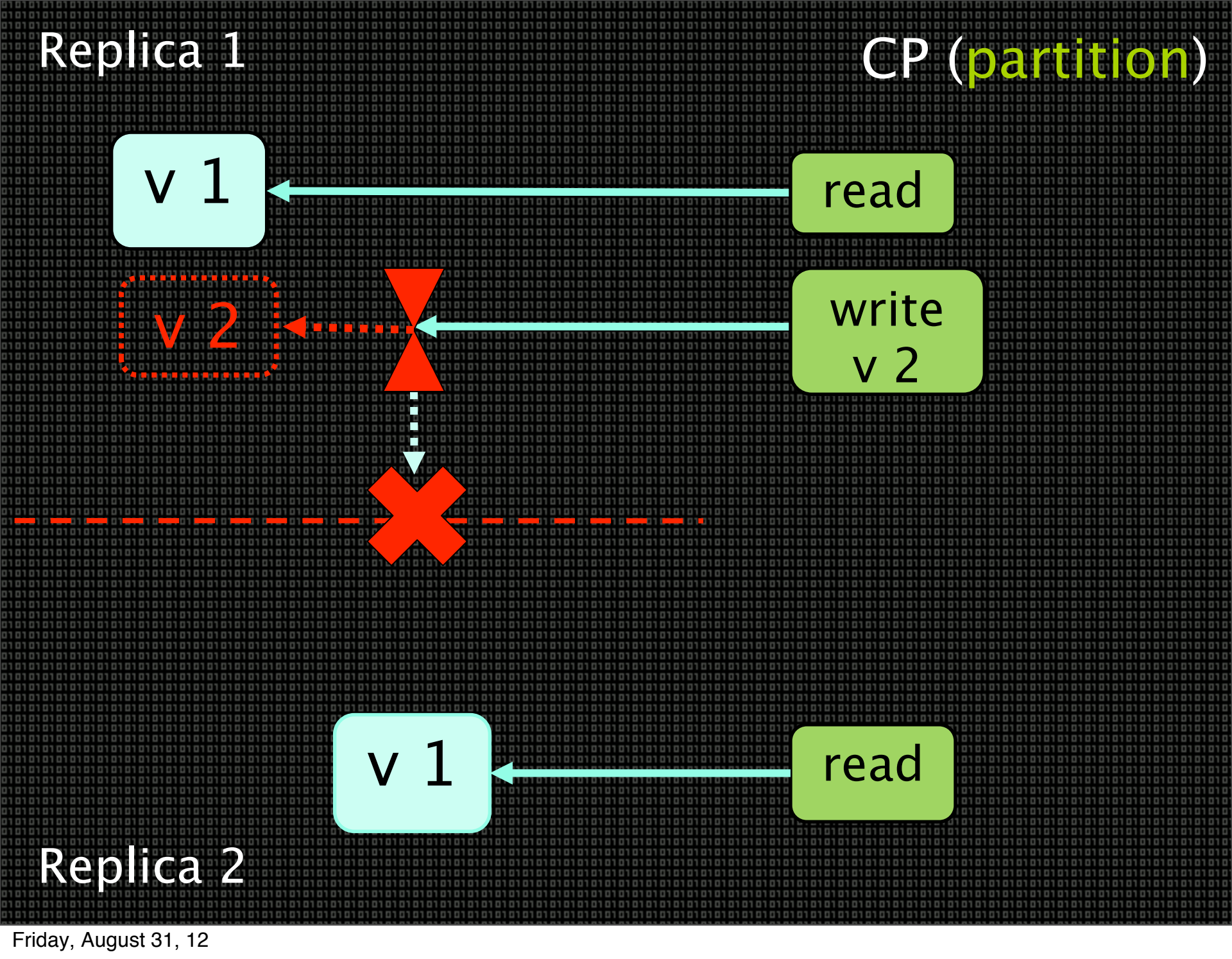
v 2

write
v 2

v 1

read

Replica 2



Replica 1

AP

v 1

v 2

write
v 2

read

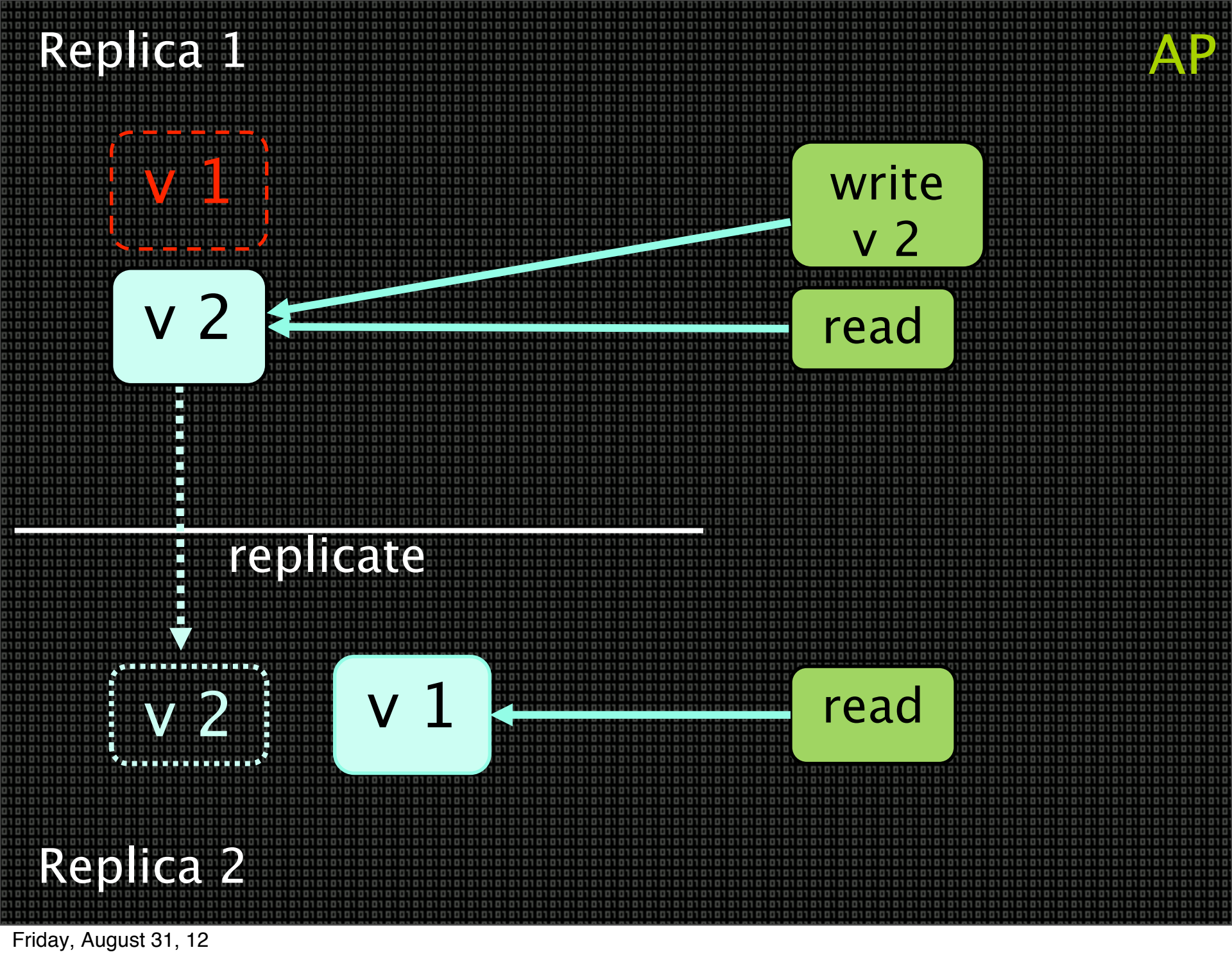
replicate

v 2

v 1

read

Replica 2



Replica 1

AP (partition)

v 1

v 2

write
v 2

read

hint
handoff

v 2

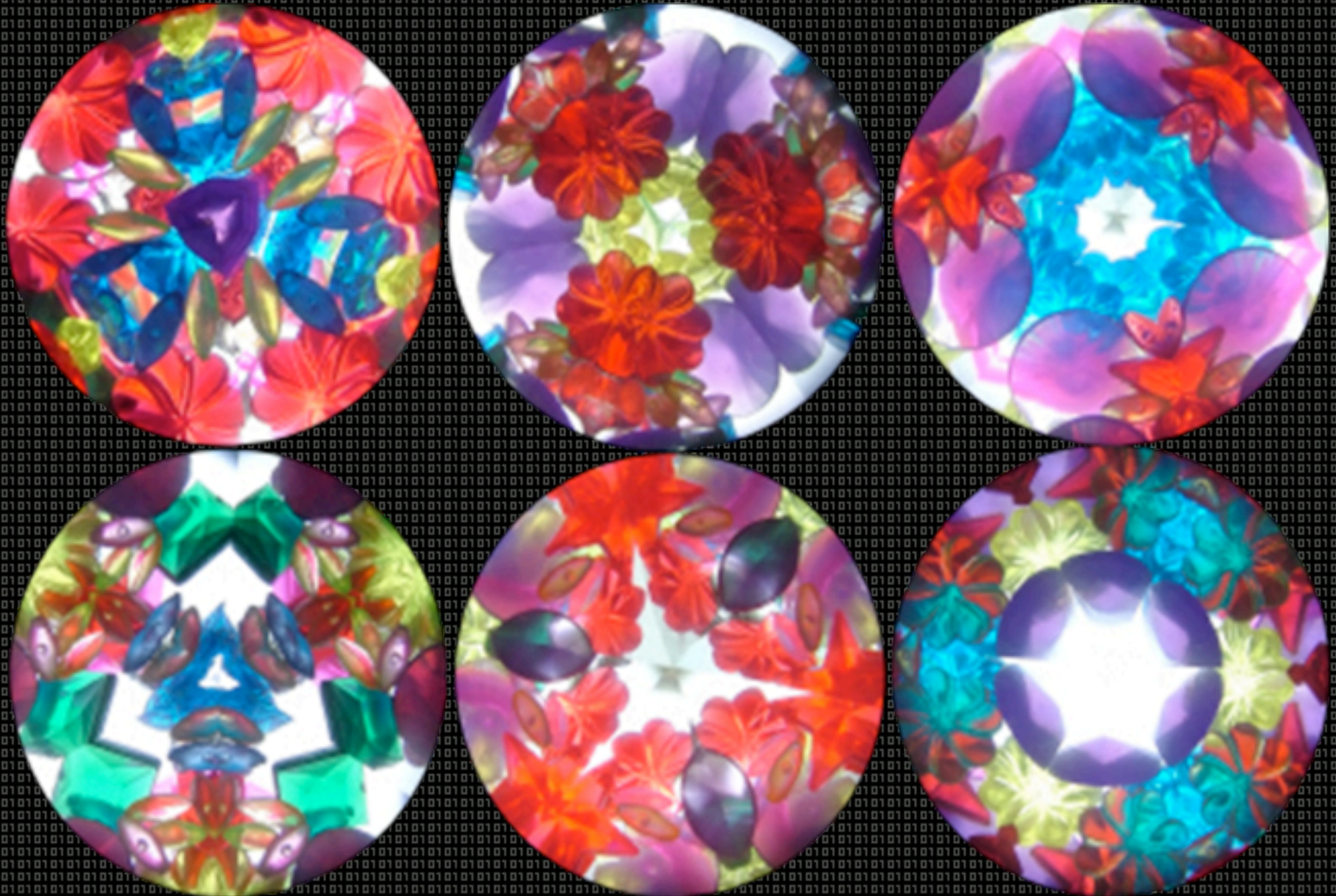


v 1

read

Replica 2

Frequent structure changes



Thank you



Many graphics I've
created myself

Some images originate from
[istockphoto.com](https://www.istockphoto.com)

except few ones taken
from Wikipedia
and product pages